

# Stochastic Separable Mixed-Integer Nonlinear Programming via Nonconvex Generalized Benders Decomposition

**Xiang Li**

Process Systems Engineering Laboratory  
Department of Chemical Engineering  
Massachusetts Institute of Technology

[xiangli@mit.edu](mailto:xiangli@mit.edu)

[xiang.li@chee.queensu.ca](mailto:xiang.li@chee.queensu.ca)

**PhD Winter School 2011**  
**Oppdal, Norway**  
**March 10, 2011**



# Agenda

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- ◆ Motivation
- ◆ Duality Theory – A Geometric Perspective
- ◆ Generalized Benders Decomposition
- ◆ Nonconvex Generalized Benders Decomposition
- ◆ Computational Study
- ◆ Summary

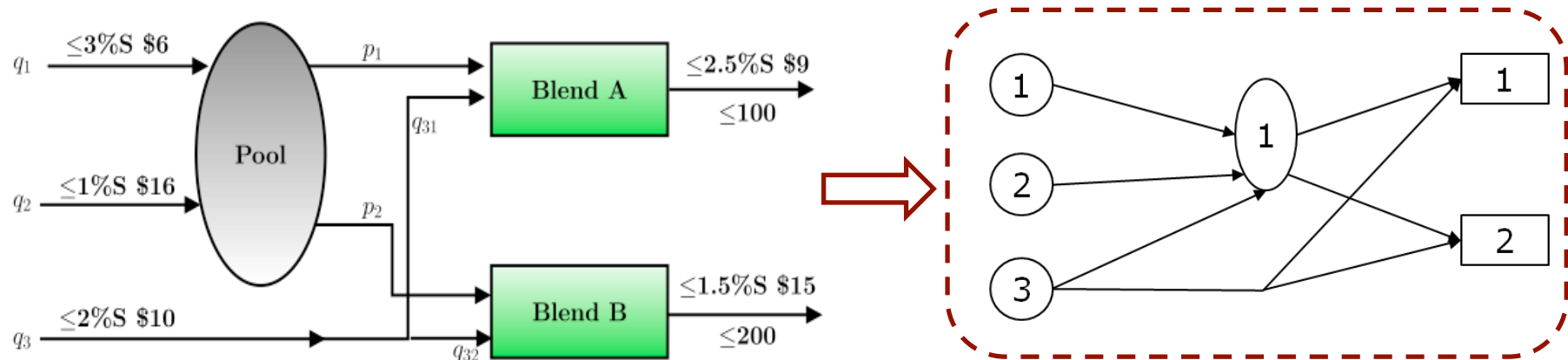
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# Motivation

## Haverly pooling problem



Haverly pooling problem \*

$$\max \quad 9(p_1 + q_{31}) + 15(p_2 + q_{32}) - (6q_1 + 16q_2 + 10q_3)$$

$$\text{s.t.} \quad q_1 + q_2 = p_1 + p_2$$

$$3q_1 + q_2 = x_s(p_1 + p_2)$$

$$q_{31} + q_{32} = q_3$$

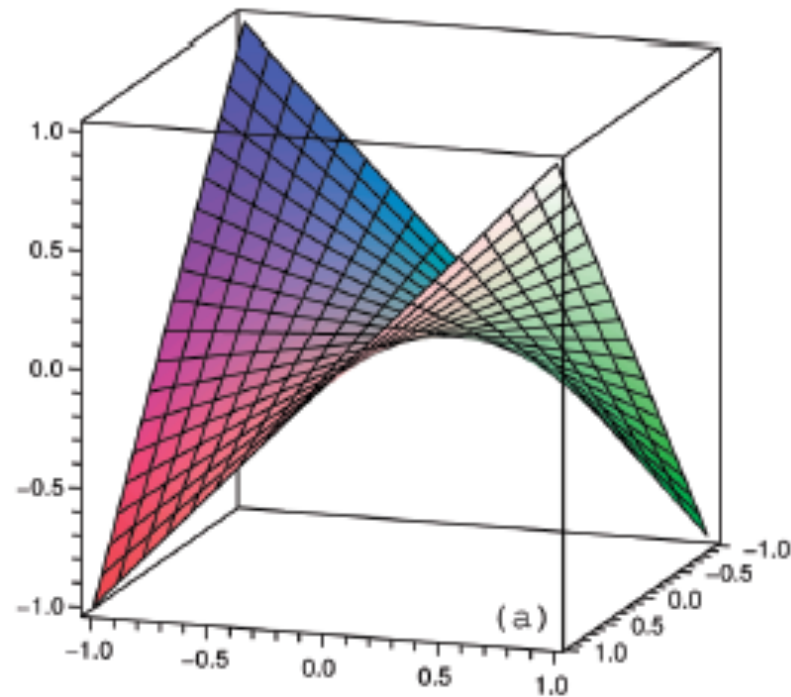
Linear inequalities

\* Haverly C. A., "Studies of the behaviour of recursion for the pooling problem", *ACM SIGMAP Bulletin*, 25:29-32, 1978.

# Motivation

## *Nonconvexity of bilinear function*

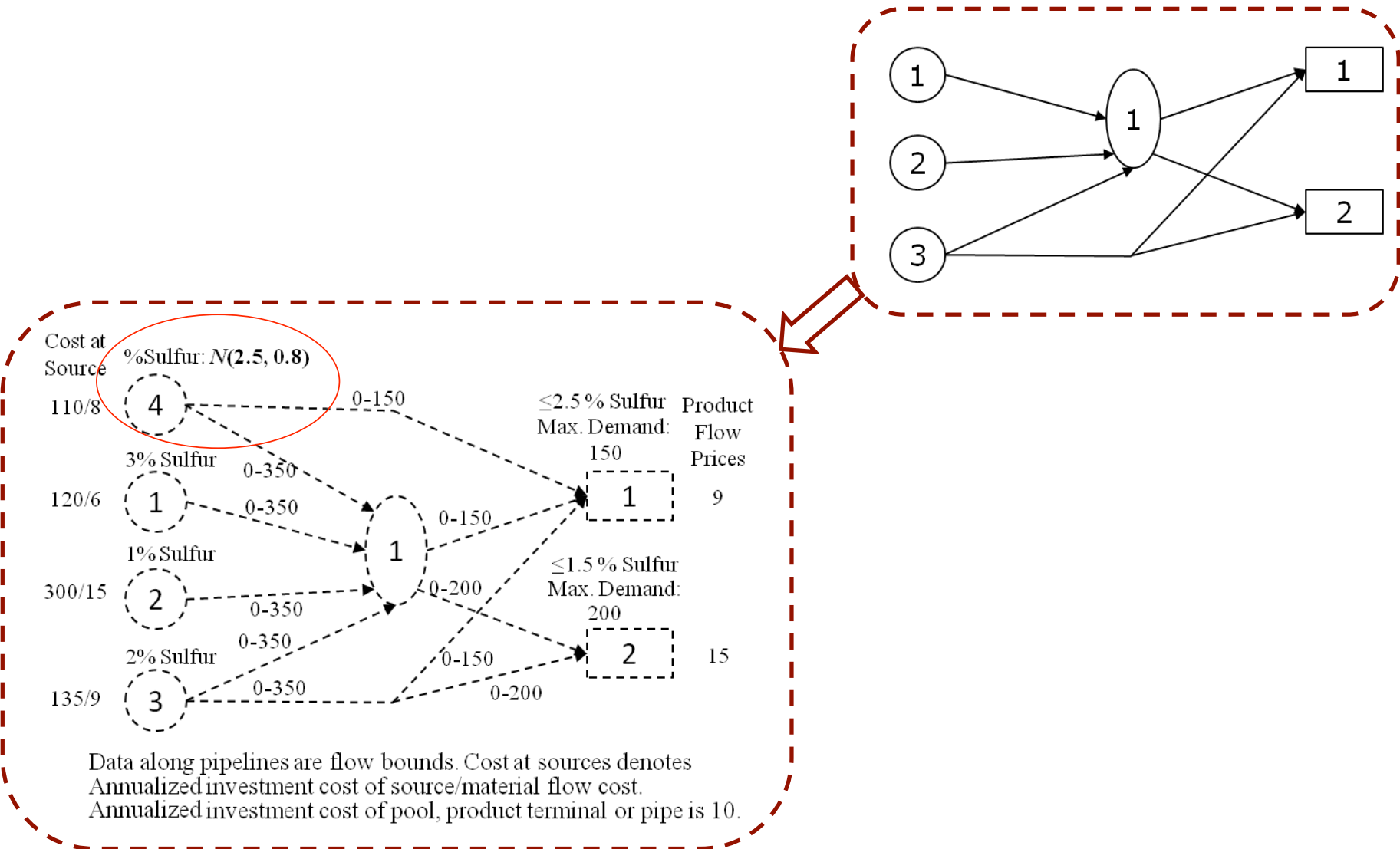
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**Bilinear function**

# Motivation

## Stochastic Haverly pooling problem



# Motivation

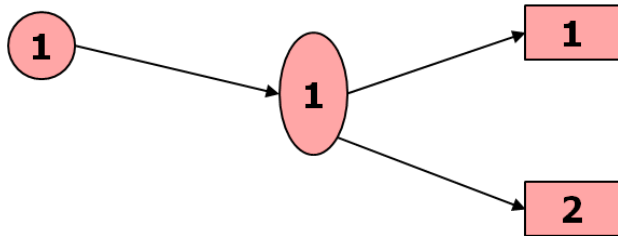
## Stochastic Haverly pooling problem – Two-stage stochastic program

**The stochastic pooling problem**  
**- A potentially large-scale mixed-integer bilinear program**

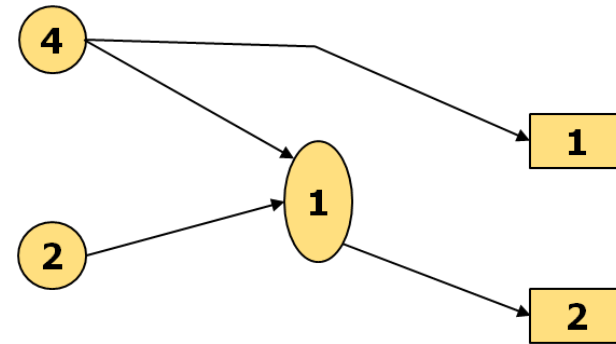
$$P = \left\{ \begin{array}{l}
 \min_{\substack{y, x_1, \dots, x_s, \\ q_1, \dots, q_s, \\ u_1, \dots, u_s}} c_1^T y + \sum_{h=1}^s (c_{2,h}^T x_h + c_{3,h}^T q_h + c_{4,h}^T u_h) \quad \left. \vphantom{\min} \right\} \text{Economic objective} \\
 \\
 s.t. \quad u_{h,l,t} = x_{h,l} q_{h,t}, \quad \forall (l,t) \in \Omega \subset \{1, \dots, n_x\} \times \{1, \dots, n_q\}, \quad \forall h \in \{1, \dots, s\} \quad \left. \vphantom{s.t.} \right\} \text{Mass balances} \\
 \quad \tilde{A}_{2,h}^{(equ)} x_h + \tilde{A}_{3,h}^{(equ)} q_h + \tilde{A}_{4,h}^{(equ)} u_h = \tilde{b}_h^{(equ)}, \quad \forall h \in \{1, \dots, s\} \\
 \\
 A_{1,h} y + A_{2,h} x_h + A_{3,h} q_h + A_{4,h} u_h \leq b_h, \quad \forall h \in \{1, \dots, s\} \quad \left. \vphantom{A_{1,h}} \right\} \text{Flow constraints} \\
 \quad \tilde{A}_{2,h} x_h + \tilde{A}_{3,h} q_h + \tilde{A}_{4,h} u_h \leq \tilde{b}_h, \quad \forall h \in \{1, \dots, s\} \\
 \quad x_h^L \leq x_h \leq x_h^U, \quad q_h^L \leq q_h \leq q_h^U, \quad \forall h \in \{1, \dots, s\} \\
 \\
 By \leq d, \quad \left. \vphantom{By} \right\} \text{Topology constraints} \\
 \\
 x_h \in \mathbb{R}^{n_x}, q_h \in \mathbb{R}^{n_q}, u_h \in \mathbb{R}^{n_u}, \quad \forall h \in \{1, \dots, s\}, y \in \{0, 1\}^{n_y}
 \end{array} \right.$$

# Motivation

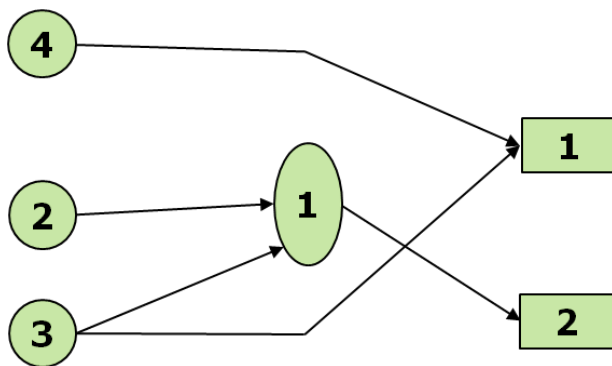
## *Stochastic Haverly pooling problem – Different Results*



**Form. 1 – Uncertainty and quality not addressed**



**Form. 2 – Only quality addressed**



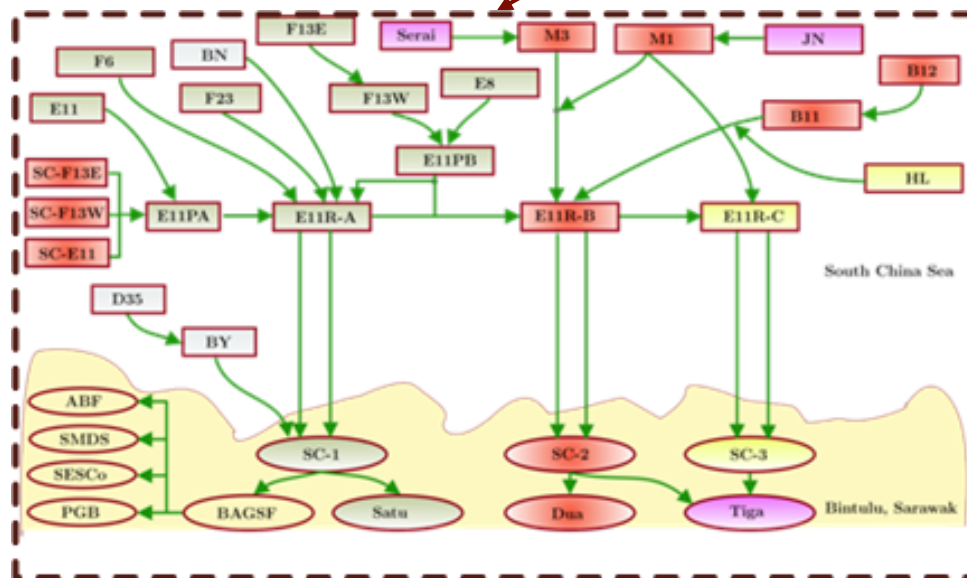
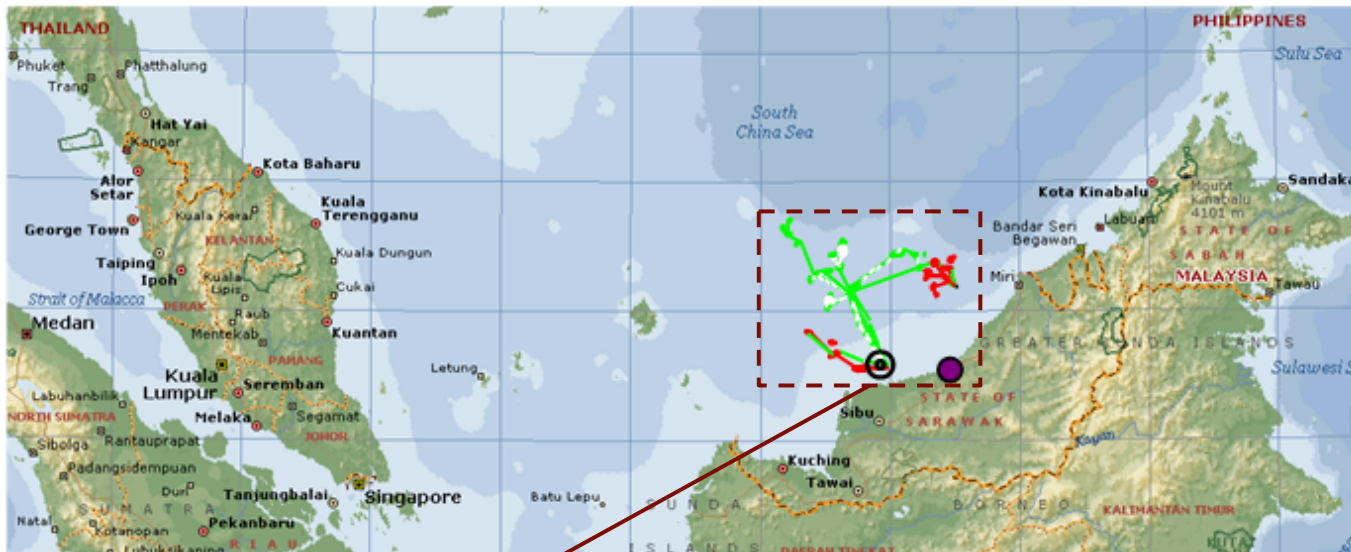
**Form. 3 – Both uncertainty and quality addressed**

Form.	$E(\text{profit})$	Sulfur Spec. Satisfied?
1	N/A	No
2	104	Yes
3	118	Yes



# Motivation

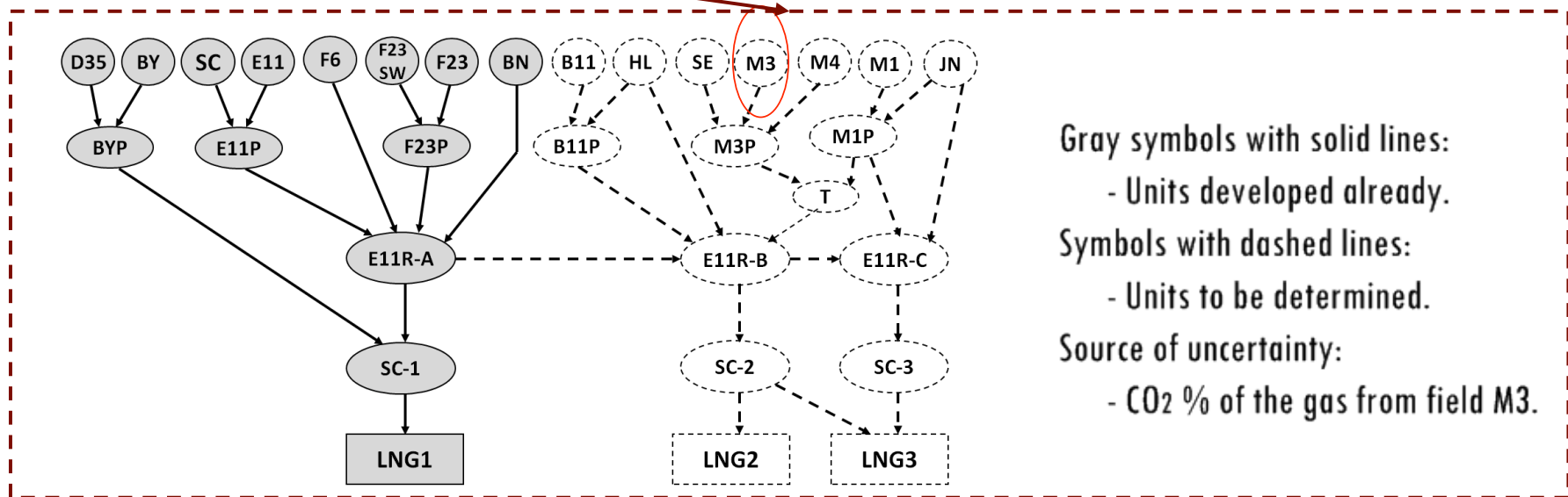
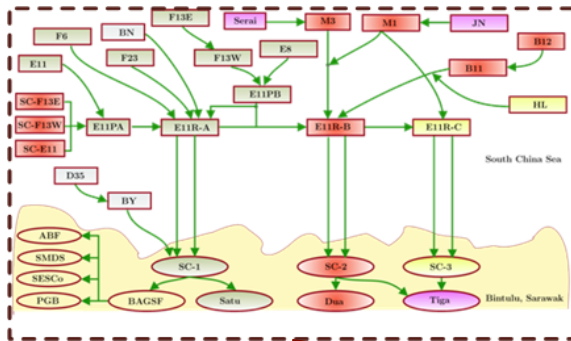
## Sarawak gas production system (SGPS)



- ◆ Daily production - 4 billion scf
- ◆ Annual revenue - US \$5 billion (4% of Malaysia's GDP)

# Motivation

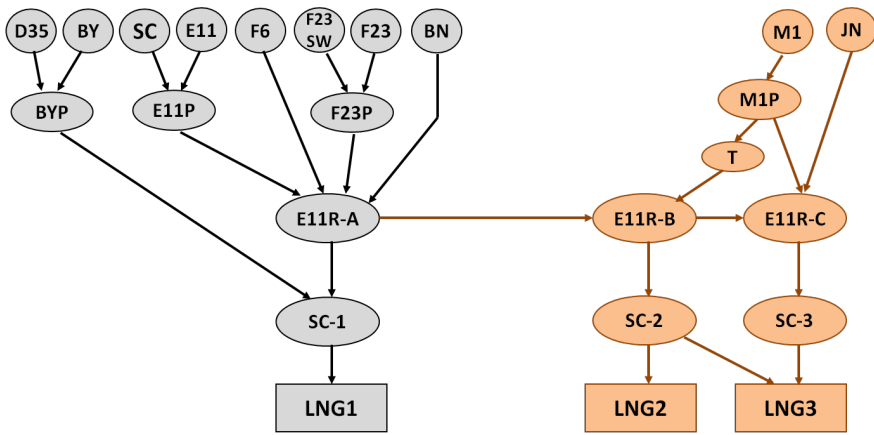
## SGPS design problem



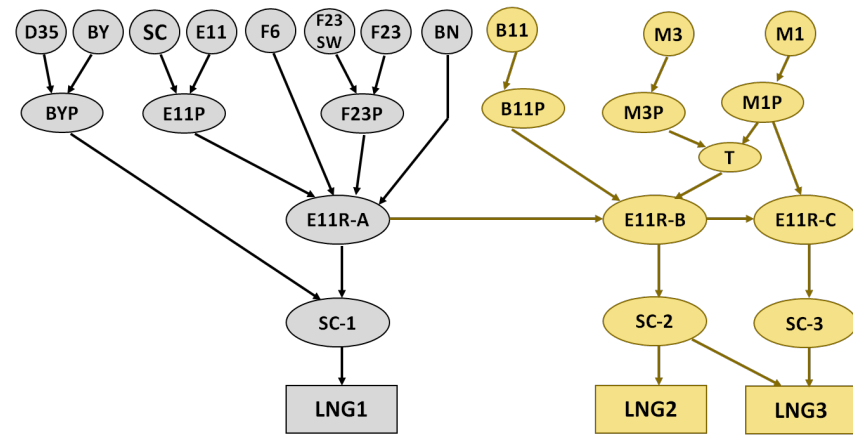
The design problem can be modeled as the stochastic pooling problem

# Motivation

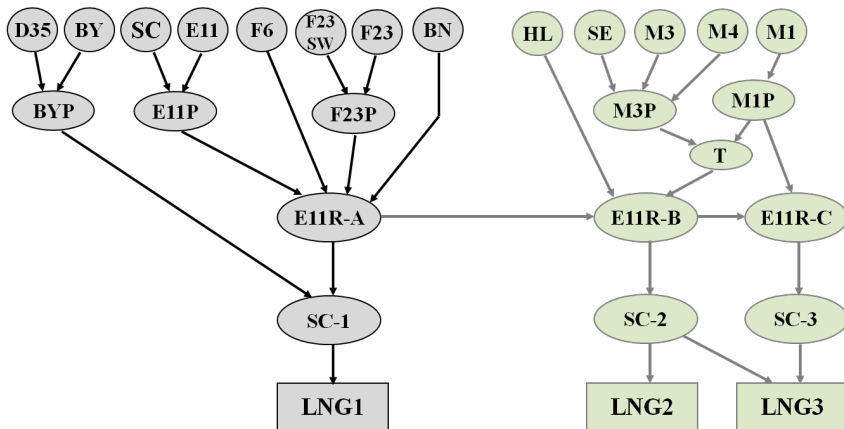
## SGPS design problem – Different results



Form. 1 – Uncertainty and quality not addressed



Form. 2 – Only quality addressed



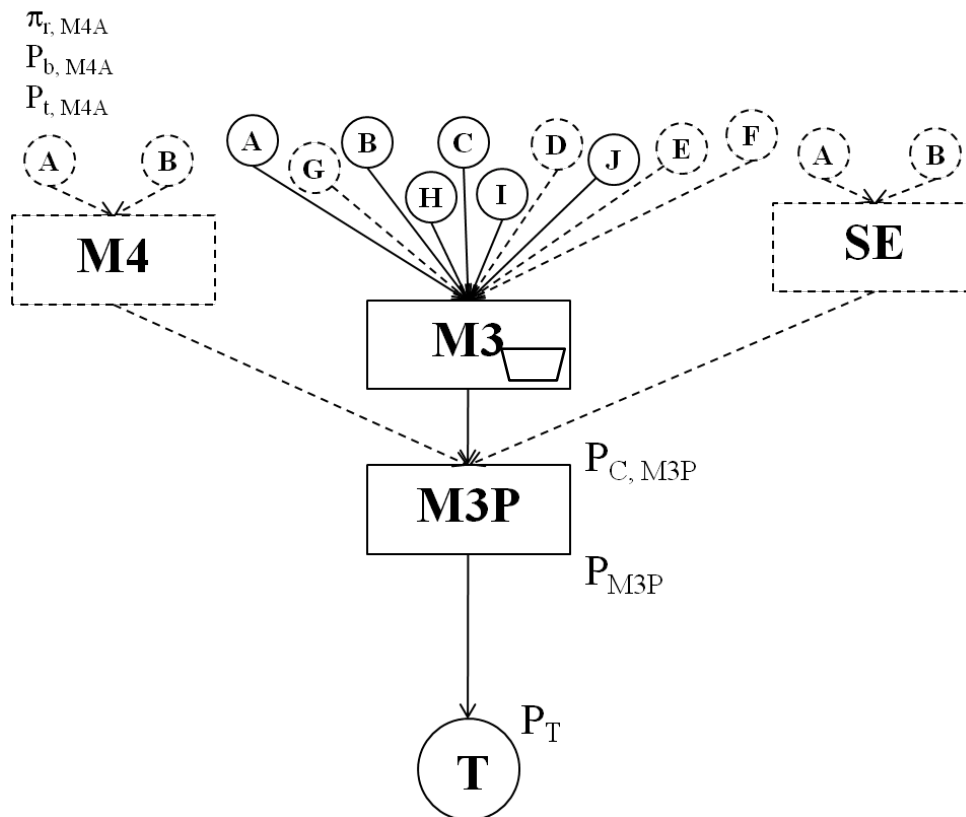
Form. 3 – Both uncertainty and quality addressed

Form.	$E(NPV)$ (Billion \$)	CO <sub>2</sub> Spec. Satisfied?
1	N/A	No
2	29.0	Yes
3	32.2	Yes

# Motivation

## SGPS design problem – More sophisticated model

### A SGPS Subsystem



#### Well performance model (M4A)

$$\alpha_{M4A} F_{M4A, M4} + \beta_{M4A} F_{M4A, M4}^2 = \pi_{r, M4A}^2 - P_{b, M4A}^2$$

$$\vartheta_{M4A} F_{M4A, M4}^2 = P_{b, M4A}^2 - \lambda_{M4A} P_{t, M4A}^2$$

#### Compression model

$$W_{M3P} = \omega_{M3P} F_{M3P, T} \left[ \left( \frac{P_{M3P}}{P_{c, M3P}} \right)^v - 1 \right]$$

#### Trunkline flow-pressure relationship

$$P_{M3P}^2 - P_T^2 = \kappa_{M3P, T} F_{M3P, T}^2$$

# Motivation

## Stochastic separable mixed-integer nonlinear programs

**Stochastic separable mixed-integer nonlinear programs**  
**– Models for energy infrastructure design and more**



**Natural Gas  
Production System**

$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y)$$

$$s.t. \quad g_h(x_h) + B_h y \leq 0, \quad h = 1, \dots, s$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

$$y \in Y$$



**Energy Polygeneration  
Plant**



**Pump Network**



**Oil Refinery**



**Software System**

# Motivation

## *Challenges for numerical optimization*

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**Stochastic separable mixed-integer nonlinear programs**  
– **Models for energy infrastructure design and more**

$$\begin{aligned} \min_{x_1, \dots, x_s, y} \quad & \sum_{h=1}^s w_h \left( f_h(x_h) + c_h^T y \right) \\ \text{s.t.} \quad & g_h(x_h) + B_h y \leq 0, \quad h = 1, \dots, s \\ & x_h \in X_h, \quad h = 1, \dots, s \\ & y \in Y \end{aligned}$$

- ◆ Accuracy of uncertainty representation
  - Large-scale optimization problems
- ◆ Profitability and feasibility of design and operation
  - Global optimization of nonconvex problems

# Motivation

## Two decomposition philosophies for scenario-based problems

$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y)$$

$$s.t. \quad g_h(x_h) + B_h y \leq 0, \quad h = 1, \dots, s$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

$$y \in Y$$

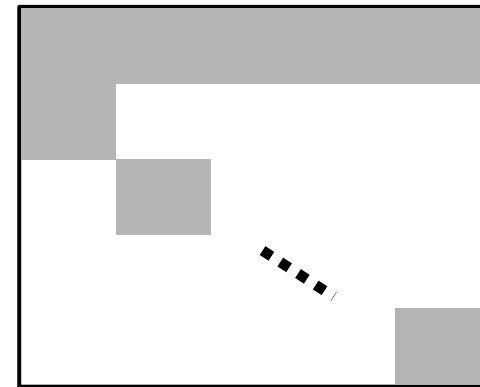
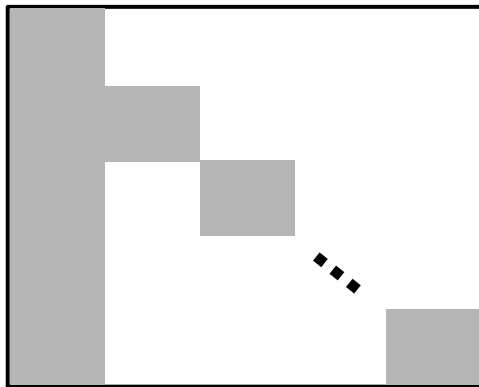
$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y_h)$$

$$s.t. \quad g_h(x_h) + B_h y_h \leq 0, \quad h = 1, \dots, s$$

$$y_h = y_{h+1}, \quad h = 1, \dots, s-1$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

$$y_h \in Y$$



- ◆ Benders decomposition/L-shaped method
  - Linear duality
- ◆ Generalized benders decomposition
  - Nonlinear duality

- ◆ Danzig-Wolfe decomposition
  - Linear duality
- ◆ Lagrangian relaxation
  - Nonlinear duality

# Motivation

## Two decomposition philosophies for scenario-based problems

$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y)$$

$$s.t. \quad g_h(x_h) + B_h y \leq 0, \quad h = 1, \dots, s$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

$$y \in Y$$

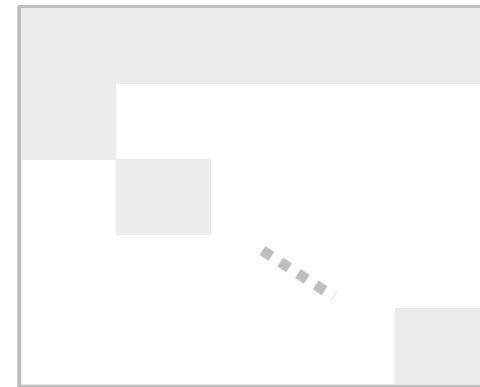
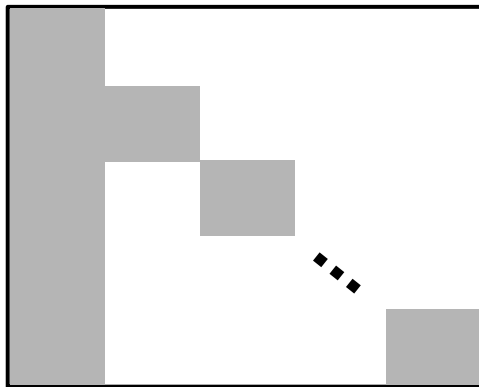
$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y_h)$$

$$s.t. \quad g_h(x_h) + B_h y_h \leq 0, \quad h = 1, \dots, s$$

$$y_h = y_{h+1}, \quad h = 1, \dots, s$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

$$y_h \in Y$$



- ◆ Benders decomposition/L-shaped method
  - Linear duality
- ◆ Generalized Benders decomposition
  - Nonlinear duality

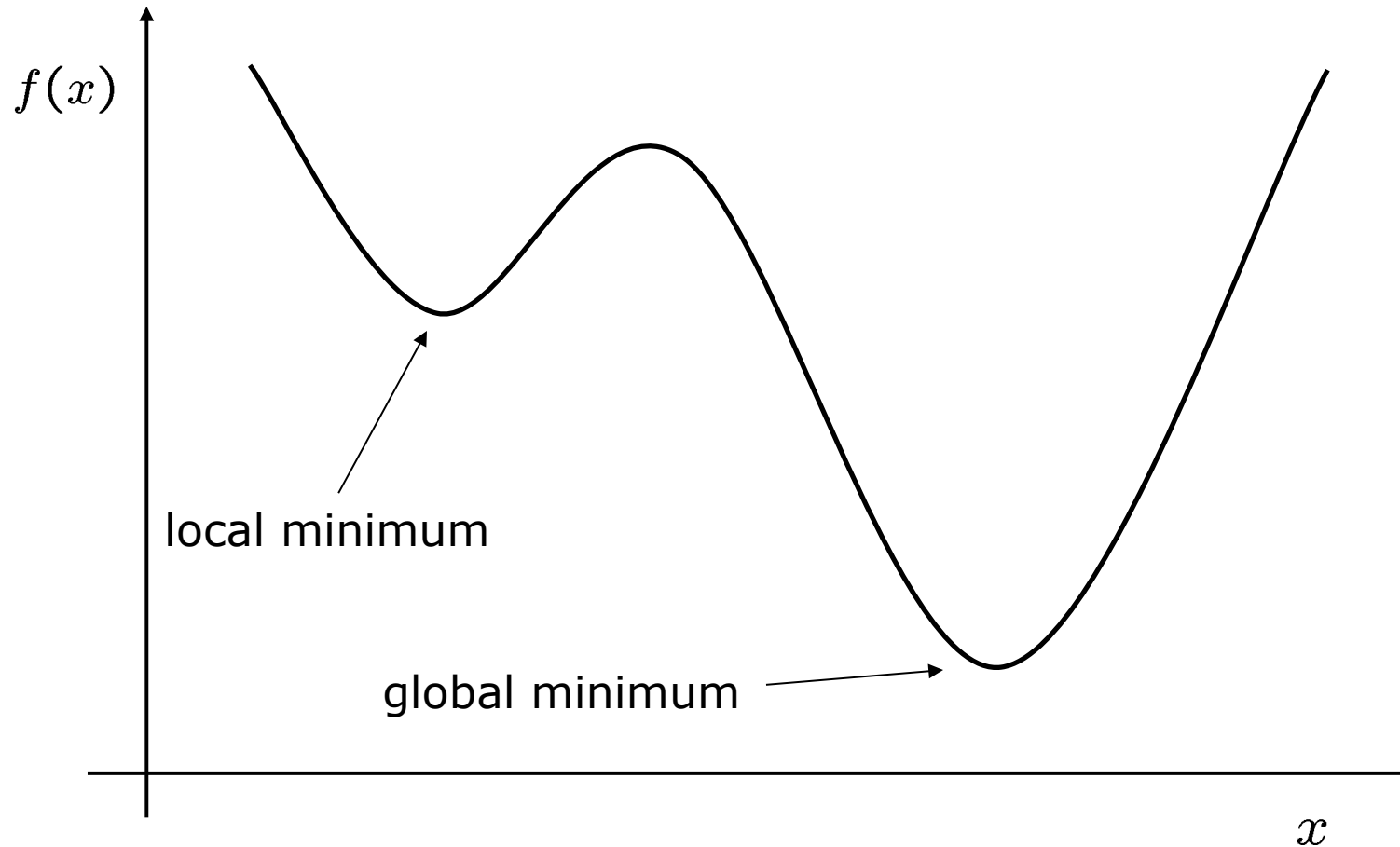
- ◆ Danzig-Wolfe decomposition
  - Linear duality
- ◆ Lagrangian relaxation
  - Nonlinear duality



# Motivation

## *Nonconvex optimization*

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Standard optimization techniques cannot distinguish between suboptimal local minima

# Motivation

## Deterministic global optimization methods for nonconvex programs

- ◆ Deterministic global optimization via branch-and-bound
  - Generates a sequence of upper and lower bounds that converge to a global optimum by domain partitioning and restriction/relaxation.
- ◆ References
  - M. Tawarmalani and N. Sahinidis. *Convexification and global optimization in continuous and mixed-integer nonlinear programming*. Kluwer Academic Publishers, 2002.
  - R. Horst, P. M. Pardalos, and N. V. Thoai. *Introduction to Global Optimization*. Kluwer Academic Publishers, 2nd edition, 2000.
  - Floudas, C. A. *Deterministic Global Optimization: Theory, Methods and Applications*; Kluwer Academic Publishers, 2000.
- ◆ Commercial software
  - BARON (Branch-and-Reduce Optimization Navigator): Sahinidis, N. V. and M. Tawarmalani, BARON 9.0.4: Global Optimization of Mixed-Integer Nonlinear Programs, User's manual, 2010.

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- ◆ Duality Theory – A Geometric Perspective
- ◆ Generalized Benders Decomposition
- ◆ Nonconvex Generalized Benders Decomposition
- ◆ Computational Study
- ◆ Summary

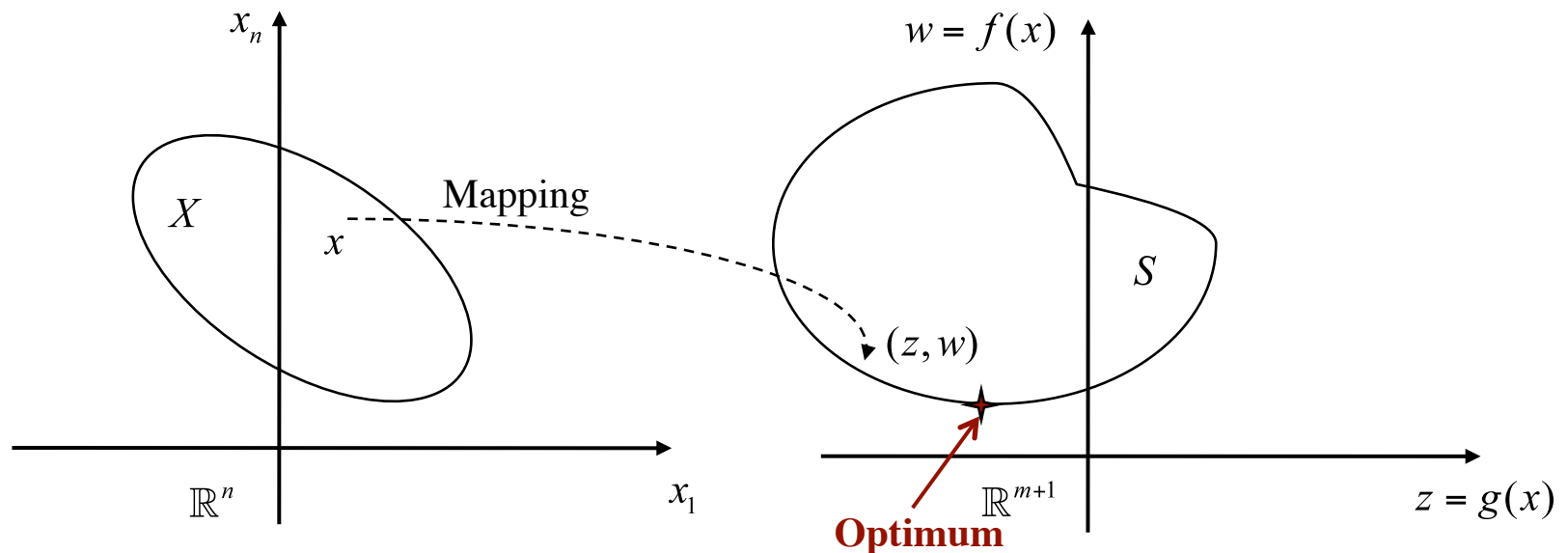
# Duality Theory – A Geometric Perspective

## *Mapping into constraint-objective space*

### Primal Problem

$$\min_{x \in X} f(x)$$

$$s.t. \quad g(x) \leq 0 \quad f: \mathbb{R}^n \rightarrow \mathbb{R}, \quad g: \mathbb{R}^n \rightarrow \mathbb{R}^m$$

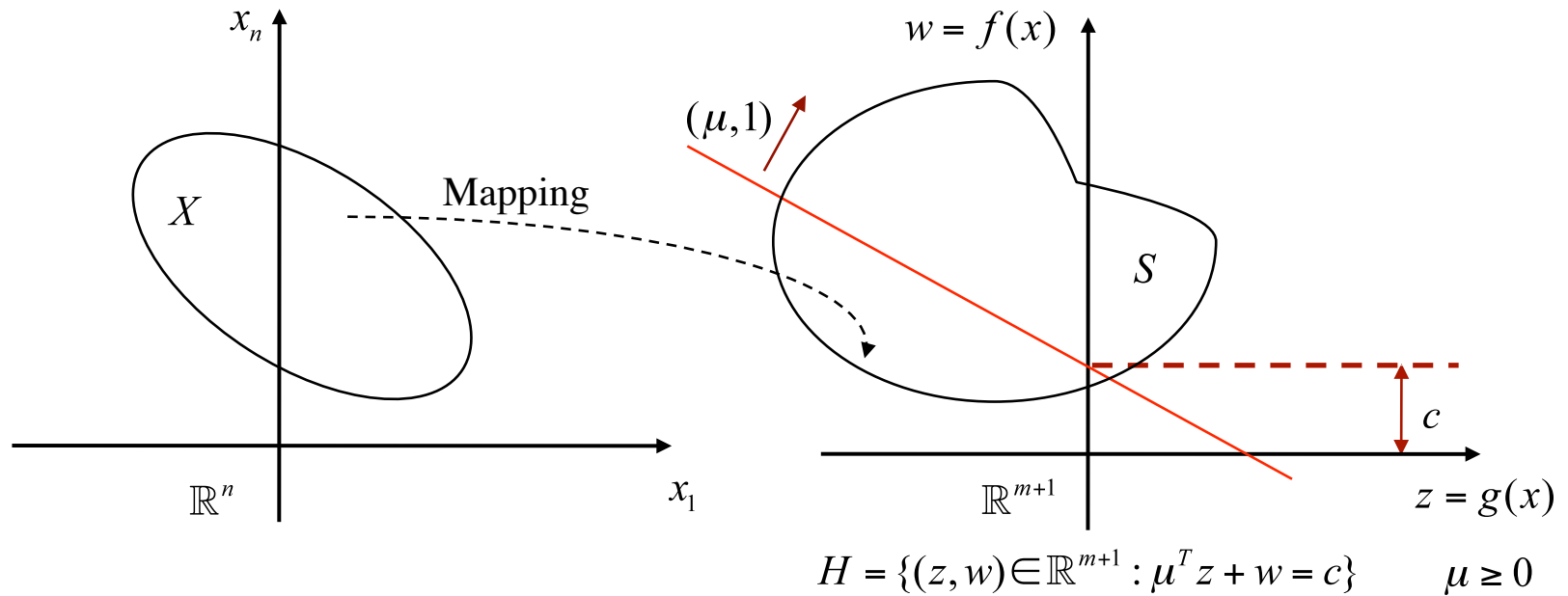


# Duality Theory – A Geometric Perspective

## Supporting hyperplane and dual function

### Primal Problem

$$\begin{aligned} \min_{x \in X} & f(x) \\ \text{s.t.} & g(x) \leq 0 \end{aligned}$$

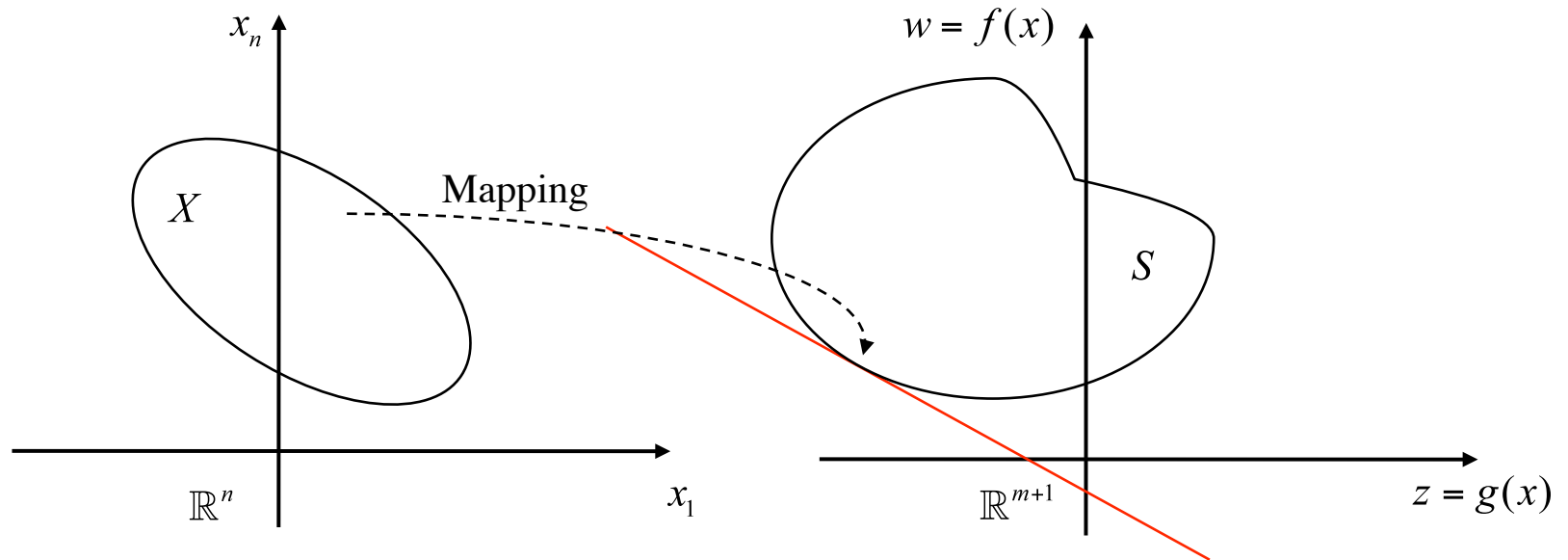


# Duality Theory – A Geometric Perspective

## Supporting hyperplane and dual function

### Primal Problem

$$\begin{aligned} \min_{x \in X} & f(x) \\ \text{s.t.} & g(x) \leq 0 \end{aligned}$$



$$\hat{H} = \{(z, w) \in \mathbb{R}^{m+1} : \mu^T z + w = \hat{c}\} \quad \mu \geq 0$$

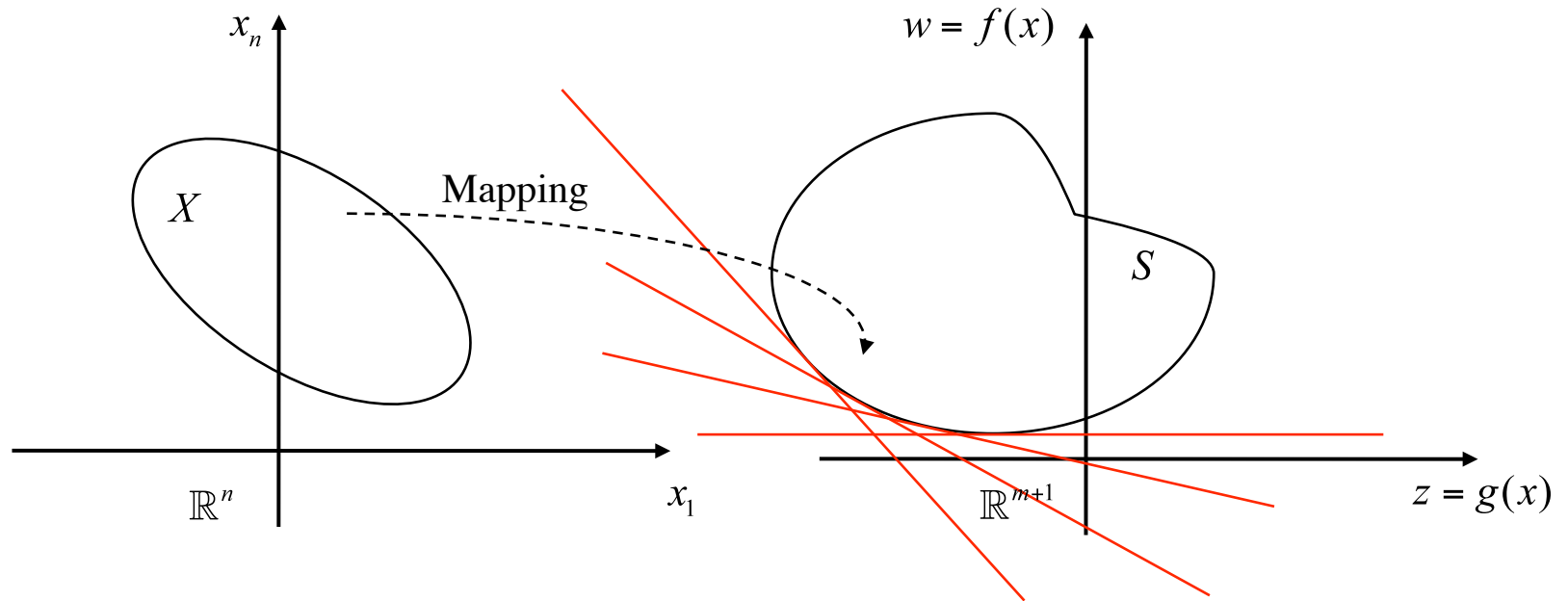
$$\hat{c} = \inf_{(z, w) \in S} \mu^T z + w = \inf_{x \in X} f(x) + \mu^T g(x) = q(\mu)$$

Dual function

# Duality Theory – A Geometric Perspective

## Dual problem

Primal Problem	Dual Problem
$\min_{x \in X} f(x)$	$\max_{\mu} q(\mu) = \inf_{x \in X} f(x) + \mu^T g(x)$
$s.t. \ g(x) \leq 0$	$s.t. \ \mu \geq 0$



$$\hat{H} = \{(z, w) \in \mathbb{R}^{m+1} : \mu^T z + w = \hat{c}\} \quad \mu \geq 0$$

$$\hat{c} = \inf_{x \in X} f(x) + \mu^T g(x) = q(\mu)$$

# Duality Theory – A Geometric Perspective

## *Weak duality*

### Primal Problem

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } g(x) \leq 0 \end{aligned}$$

### Dual Problem

$$\begin{aligned} \max_{\mu} q(\mu) = \inf_{x \in X} f(x) + \mu^T g(x) \\ \text{s.t. } \mu \geq 0 \end{aligned}$$

**Weak Duality:** Define  $f^* = \inf_{x \in X, g(x) \leq 0} f(x)$ ,  $g^* = \sup_{\mu \geq 0} q(\mu)$ , then  $f^* \geq g^*$ .

Proof: For all  $\mu \geq 0$  and for all  $x \in X$  such that  $g(x) \leq 0$ ,

$$q(\mu) = \inf_{x \in X} L(x, \mu) \leq f(x) + \mu^T g(x) \leq f(x),$$

so

$$q^* = \sup_{\mu \geq 0} q(\mu) \leq \inf_{x \in X, g(x) \leq 0} f(x) = f^*.$$



# Duality Theory – A Geometric Perspective

## *Strong duality*

### Primal Problem

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } g(x) \leq 0 \end{aligned}$$

### Dual Problem

$$\begin{aligned} \max_{\mu} q(\mu) = \inf_{x \in X} f(x) + \mu^T g(x) \\ \text{s.t. } \mu \geq 0 \end{aligned}$$

**Strong Duality:** Let  $-\infty < f^* = \inf_{x \in X, g(x) \leq 0} f(x) < +\infty$ ,  $g^* = \sup_{\mu \geq 0} q(\mu)$ . If the primal problem is convex, feasible and satisfies Slater's condition, then  $f^* = g^*$ .

Proof: See

D. P. Bertsekas. *Nonlinear Programming*. Athena Scientific, 2nd edition, 1999.

# Duality Theory – A Geometric Perspective

## Duality gap

### Primal Problem

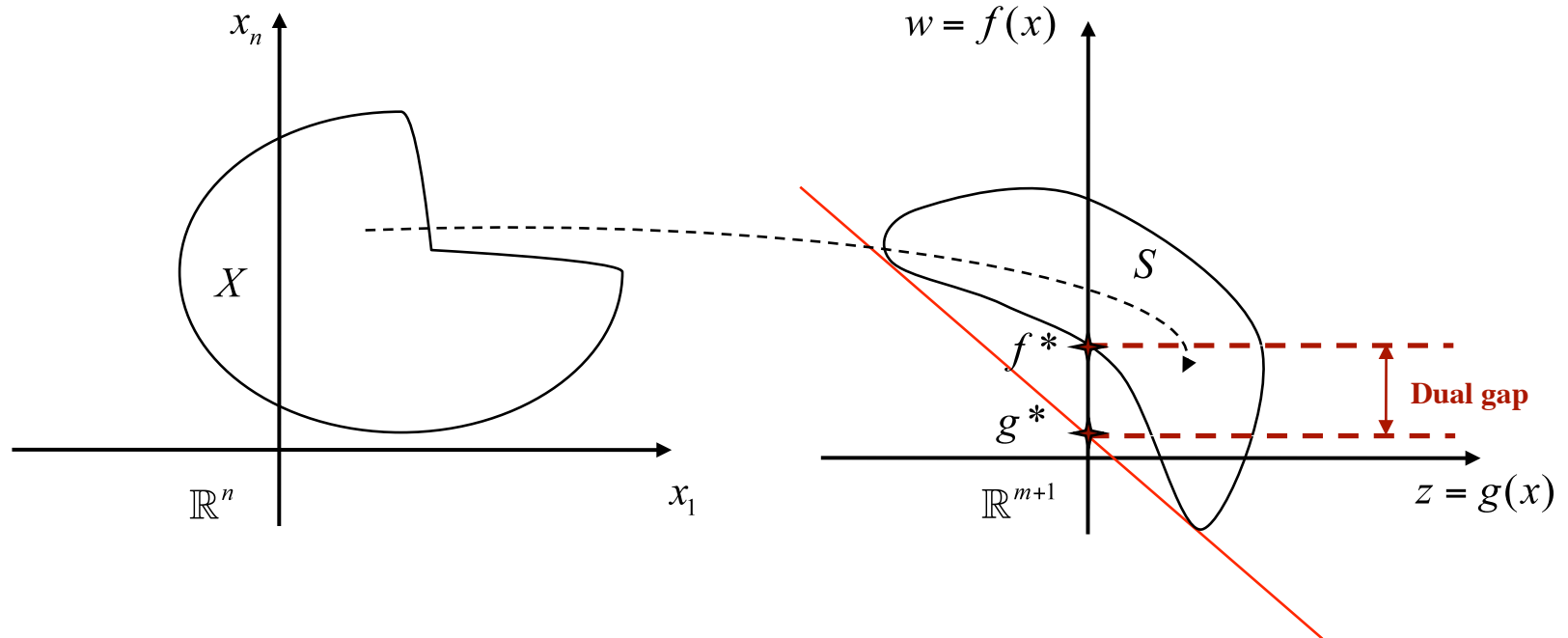
$$\min_{x \in X} f(x)$$

$$s.t. \quad g(x) \leq 0$$

### Dual Problem

$$\max_{\mu} \quad q(\mu) = \inf_{x \in X} f(x) + \mu^T g(x)$$

$$s.t. \quad \mu \geq 0$$



# Duality Theory – A Geometric Perspective

## Example 1

### Primal Problem

$$\min_{x \in X} x_2 - x_1$$

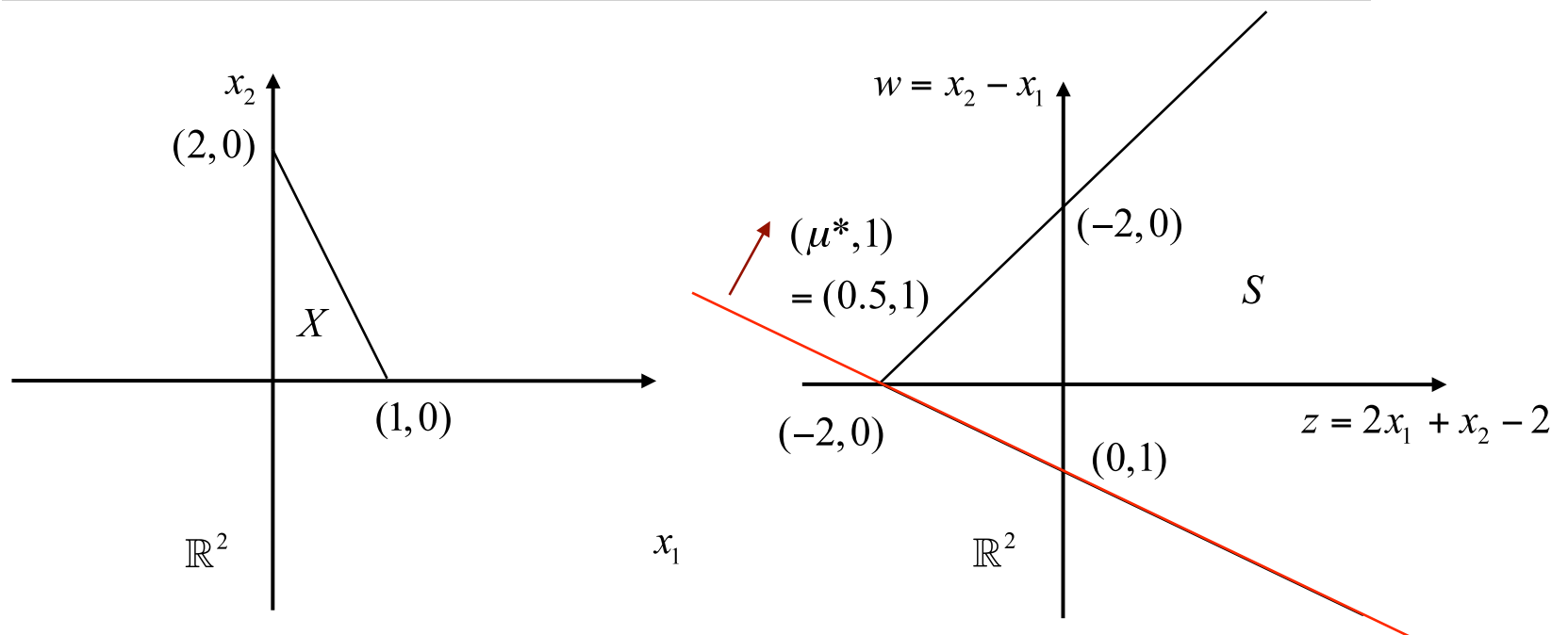
$$s.t. \quad 2x_1 + x_2 - 2 \leq 0$$

$$X = \{x \in \mathbb{R}^2 : x_1 \geq 0, x_2 \geq 0\}$$

### Dual Problem

$$\max_{\mu} \inf_{x \in X} (2\mu - 1)x_1 + (\mu + 1)x_2 - 2\mu$$

$$s.t. \quad \mu \geq 0$$



# Duality Theory – A Geometric Perspective

## Example 2

### Primal Problem

$$\min_{x \in X} (x_1 - 1)^2 + (x_2 - 1)^2$$

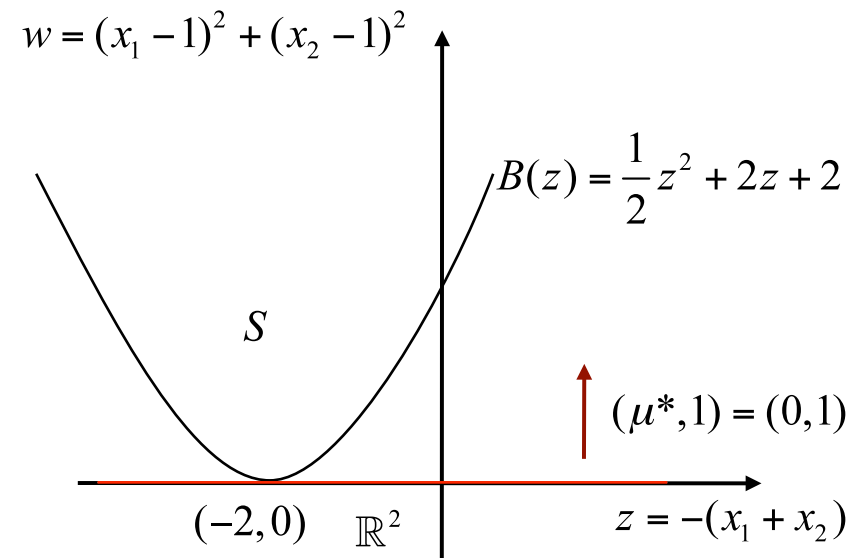
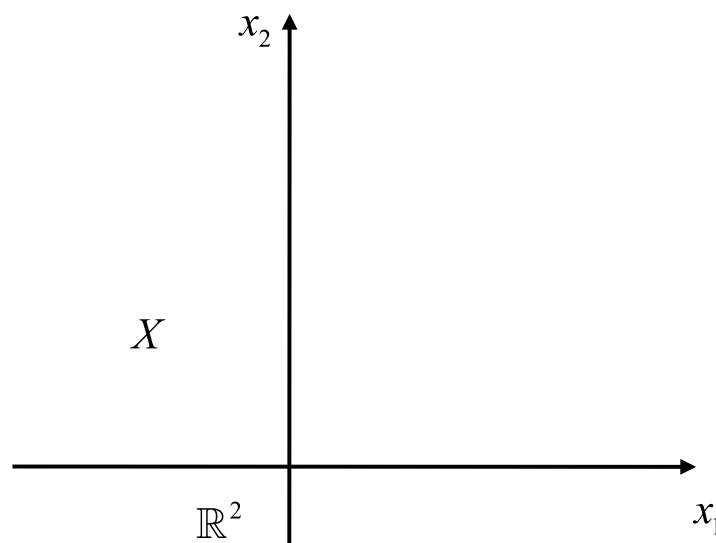
$$s.t. \quad -(x_1 + x_2) \leq 0$$

$$X = \mathbb{R}^2$$

### Dual Problem

$$\max_{\mu} \inf_{x \in \mathbb{R}^2} (x_1 - 1)^2 + (x_2 - 1)^2 - \mu(x_1 + x_2)$$

$$s.t. \quad \mu \geq 0$$



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- ◆ **Generalized Benders Decomposition**
- ◆ Nonconvex Generalized Benders Decomposition
- ◆ Computational Study
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# Generalized Benders Decomposition

## *A generalization of Benders decomposition*

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- ◆ Generalized Benders decomposition (GBD) is an extension of Benders decomposition.
  - A.M. Geoffrion. Generalized Benders decomposition. *Journal of Optimization Theory and Applications*, 10(4):237–260, 1972.
  - J. F. Benders. Partitioning procedures for solving mixed-variables programming problems. *Numerische Mathematik*, 4:238–252, 1962.

- ◆ Target Problem

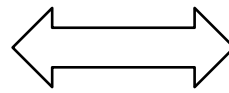
$$\begin{aligned} \min_{x,y} \quad & f(x, y) \\ \text{s.t.} \quad & g(x, y) \leq 0 \\ & x \in X, y \in Y \end{aligned}$$

- Assumptions:**
- (a)  $X \subset \mathbb{R}^{n_x}$  and  $Y \subset \mathbb{R}^{n_y}$  are compact,
  - (b) For all  $\hat{y} \in Y$ ,  $f(\cdot, y): \mathbb{R}^{n_x} \rightarrow \mathbb{R}$  and  $g(\cdot, y): \mathbb{R}^{n_x} \rightarrow \mathbb{R}^m$  are convex on  $X$ .
  - (c) For  $y$  fixed to any feasible  $\hat{y} \in Y$ , the problem satisfies Slater's condition.

# Generalized Benders Decomposition

## Principle of projection

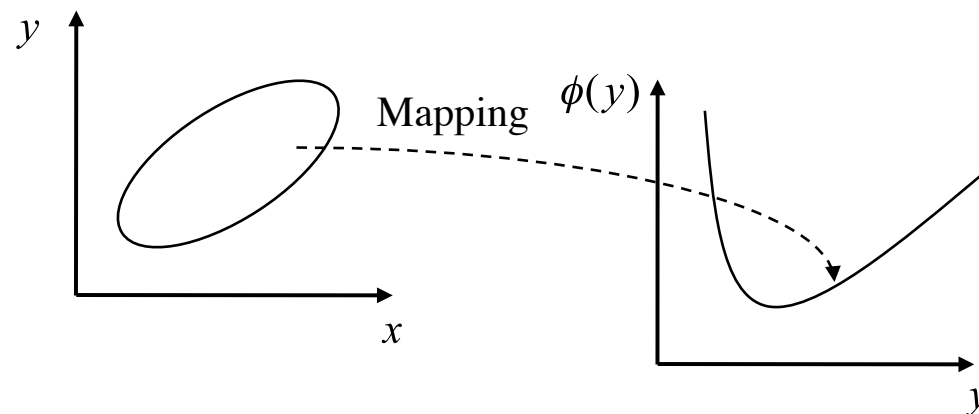
$$\begin{aligned} \min_{x,y} f(x,y) \\ \text{s.t. } g(x,y) \leq 0 \\ x \in X, y \in Y \end{aligned}$$



$$\begin{aligned} \min_y \phi(y) \\ \text{s.t. } \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y) \\ y \in Y \cap V, \\ V = \{y : \exists x \in X, g(x,y) \leq 0\} \end{aligned}$$

Optimality

Feasibility



# Generalized Benders Decomposition

## *Optimality cuts*

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$$\min_y \phi(y)$$

$$s.t. \quad \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y) \quad \text{Optimality}$$

$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x,y) \leq 0\} \quad \text{Feasibility}$$



# Generalized Benders Decomposition

## Optimality

$$\min_y \phi(y)$$

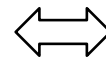
$$s.t. \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y)$$

$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x,y) \leq 0\}$$

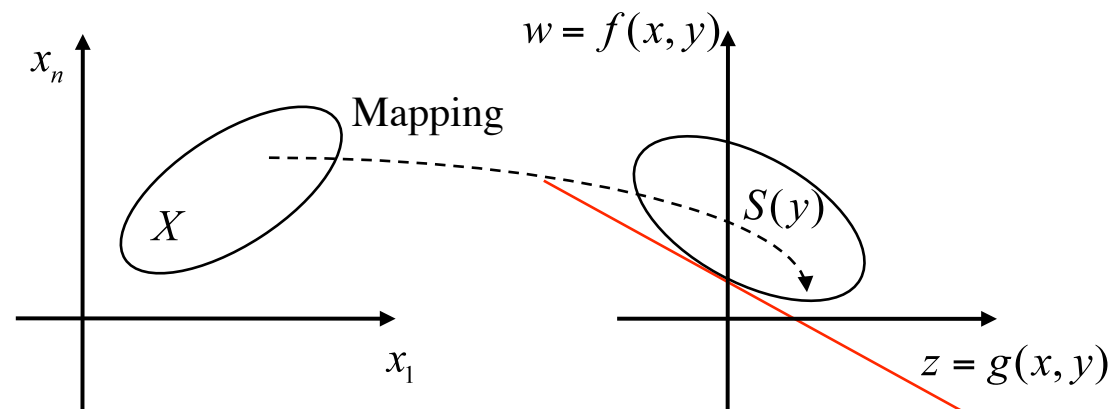
Convex, Slater's  
condition holds

Optimality



$$\phi(y) = \sup_{\lambda \geq 0} \inf_{x \in X} f(x,y) + \lambda^T g(x,y)$$

Feasibility



# Generalized Benders Decomposition

## Optimality, nonconvexity strikes

$$\min_y \phi(y)$$

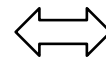
$$s.t. \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y)$$

$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x,y) \leq 0\}$$

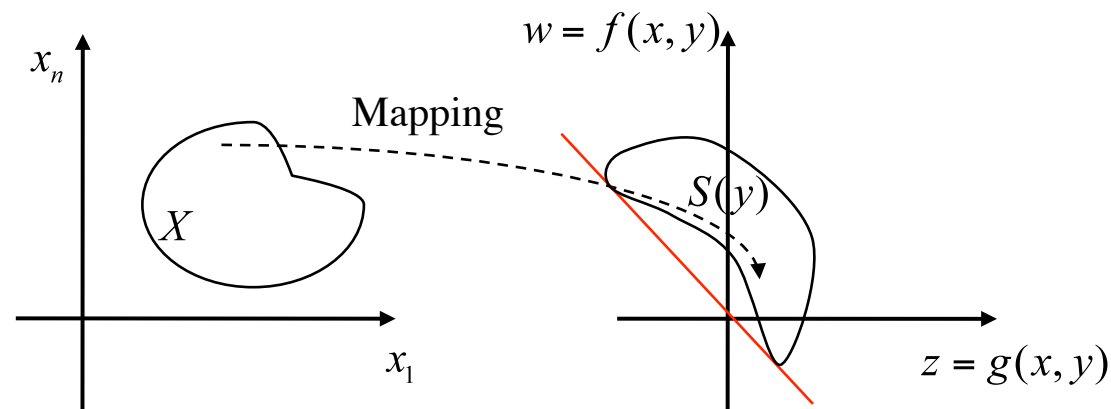
Convex, Slater's  
condition holds

Optimality



$$\phi(y) = \sup_{\lambda \geq 0} \inf_{x \in X} f(x,y) + \lambda^T g(x,y)$$

Feasibility



# Generalized Benders Decomposition

## *Feasibility*

---

$$\min_y \phi(y)$$

$$s.t. \quad \phi(y) = \inf_{x \in X, g(x, y) \leq 0} f(x, y) \quad \text{Optimality}$$

$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x, y) \leq 0\} \quad \text{Feasibility}$$

# Generalized Benders Decomposition

## Feasibility

$$\min_y \phi(y)$$

$$s.t. \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y) \quad \text{Optimality}$$

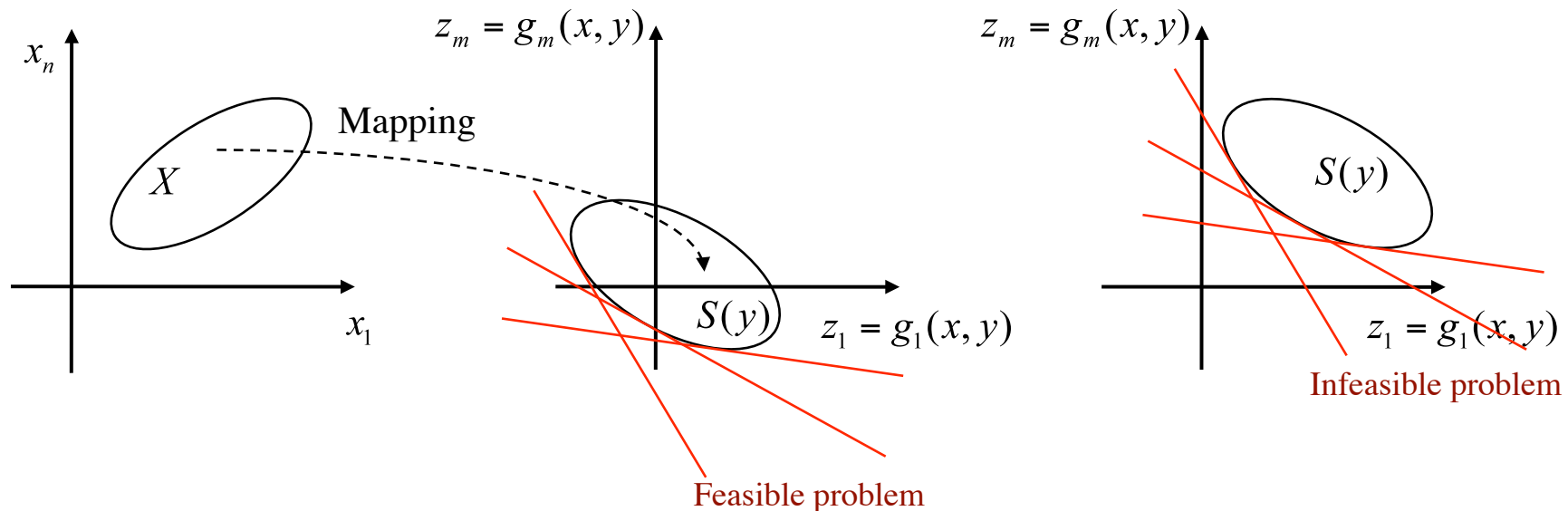
$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x,y) \leq 0\}$$

Feasibility  $\iff$  Convex

$$\{y : 0 \geq \inf_{x \in X} \mu^T g(x,y), \quad \forall \mu \in M\}$$

$$M = \{\mu : \mu \geq 0, \sum_i \mu_i = 1\}$$



# Generalized Benders Decomposition

## Feasibility cuts, nonconvexity strikes

$$\min_y \phi(y)$$

$$s.t. \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y) \quad \text{Optimality}$$

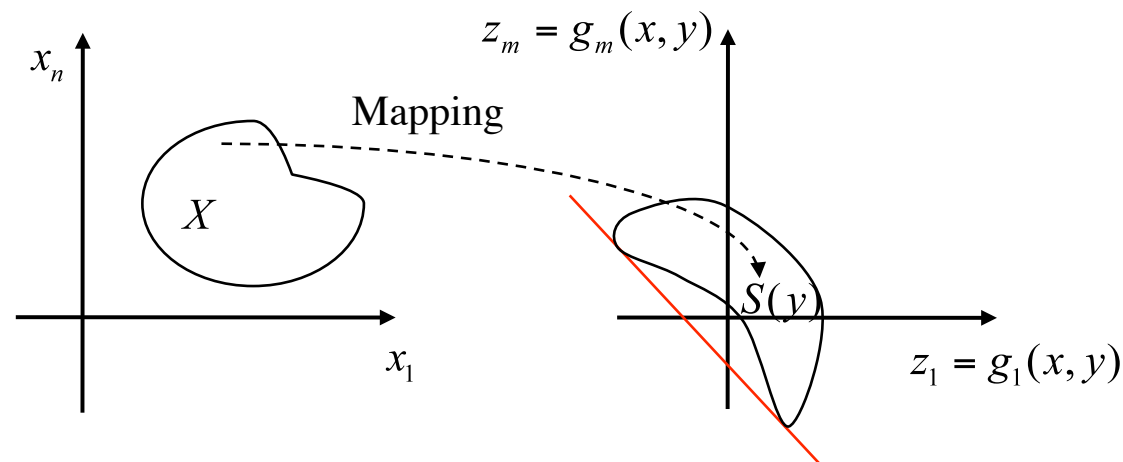
$$y \in Y \cap V,$$

$$V = \{y : \exists x \in X, g(x,y) \leq 0\}$$

Feasibility  $\longleftrightarrow$  Convex

$$\{y : 0 \geq \inf_{x \in X} \mu^T g(x,y), \quad \forall \mu \in M\}$$

$$M = \{\mu : \mu \geq 0, \sum_i \mu_i = 1\}$$



# Generalized Benders Decomposition

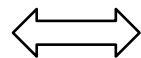
## Master problem

$$\begin{aligned} \min_{x,y} & f(x,y) \\ \text{s.t.} & g(x,y) \leq 0 \\ & x \in X, y \in Y \end{aligned}$$



$$\begin{aligned} \min_y & \phi(y) \\ \text{s.t.} & \phi(y) = \inf_{x \in X, g(x,y) \leq 0} f(x,y) \\ & y \in Y \cap V, \\ & V = \{y : \exists x \in X, g(x,y) \leq 0\} \end{aligned}$$

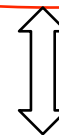
**Dualization**



**Master Problem**

$$\begin{aligned} \min_{y,\eta} & \eta \\ \text{s.t.} & \eta \geq \inf_{x \in X} f(x,y) + \lambda^T g(x,y), \quad \forall \lambda \geq 0 \\ & 0 \geq \inf_{x \in X} \mu^T g(x,y), \quad \forall \mu \in M \end{aligned}$$

Optimality cuts  
Feasibility cuts



$$\begin{aligned} \min_y & \phi(y) \\ \text{s.t.} & \phi(y) = \sup_{\lambda \geq 0} \inf_{x \in X} f(x,y) + \lambda^T g(x,y) \\ & 0 \geq \inf_{x \in X} \mu^T g(x,y), \quad \forall \mu \in M \end{aligned}$$

# Generalized Benders Decomposition

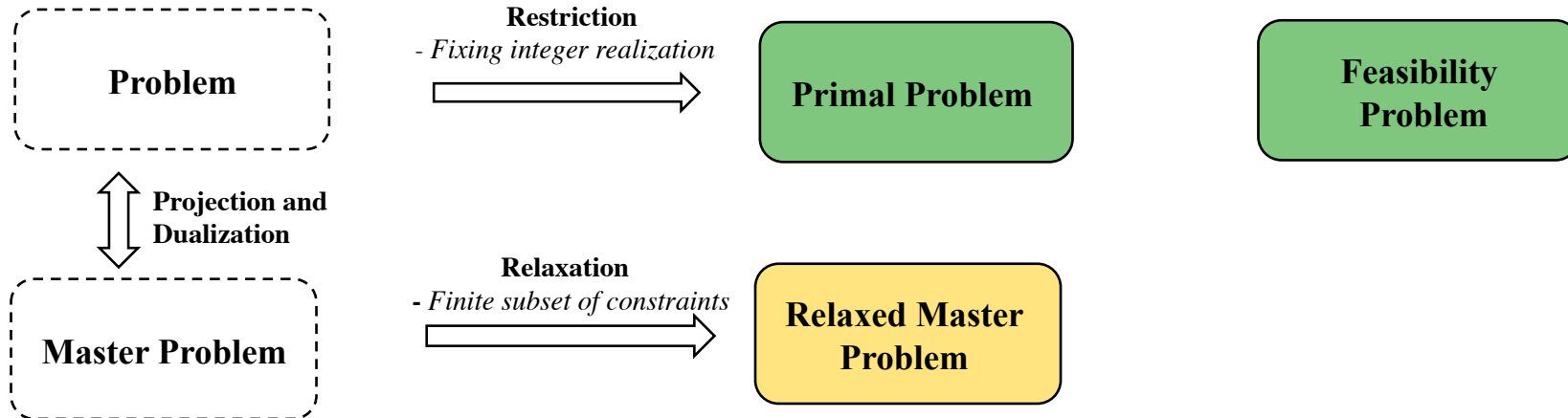
## Subproblems generated via restriction and relaxation

---

$$\begin{aligned} \min_{x,y} & f(x, y) \\ \text{s.t.} & g(x, y) \leq 0 \\ & x \in X, y \in Y \end{aligned}$$

$$\begin{aligned} \text{obj}_{\text{Primal}}^k &= \min_x f(x, y^{(k)}) \\ \text{s.t.} & g(x, y^{(k)}) \leq 0 \\ & x \in X \end{aligned}$$

$$\begin{aligned} \text{obj}_{\text{Feas}}^k &= \min_x z \\ \text{s.t.} & g(x, y^{(k)}) \leq z, \\ & x \in X, z \in Z \subset \{z \in \mathbb{R}^m : z \geq 0\} \end{aligned}$$



$$\begin{aligned} \min_{y,\eta} & \eta \\ \text{s.t.} & \eta \geq \inf_{x \in X} f(x, y) + \lambda^T g(x, y), \\ & \forall \lambda \geq 0 \\ & 0 \geq \inf_{x \in X} \mu^T g(x, y), \\ & \forall \mu \in M \end{aligned}$$

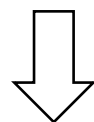
$$\begin{aligned} \min_{y,\eta} & \eta \\ \text{s.t.} & \eta \geq \inf_{x \in X} f(x, y) + (\lambda^j)^T g(x, y), \quad \forall \lambda^j \in T^k \\ & 0 \geq \inf_{x \in X} (\mu^i)^T g(x, y), \quad \forall \mu^i \in S^k \end{aligned}$$

# Generalized Benders Decomposition

## Relaxed master problem with separability in $x$ and $y$

### Relaxed Master Problem

$$\begin{aligned} \min_{y, \eta} \quad & \eta \\ \text{s.t.} \quad & \eta \geq \inf_{x \in X} f(x, y) + (\lambda^j)^T g(x, y), \quad \forall \lambda^j \in T^k \\ & 0 \geq \inf_{x \in X} (\mu^i)^T g(x, y), \quad \forall \mu^i \in S^k \end{aligned}$$



$$f(x, y) \equiv f_1(x) + f_2(y), \quad g(x, y) \equiv g_1(x) + g_2(y) *$$

This is implied in Benders decomposition!

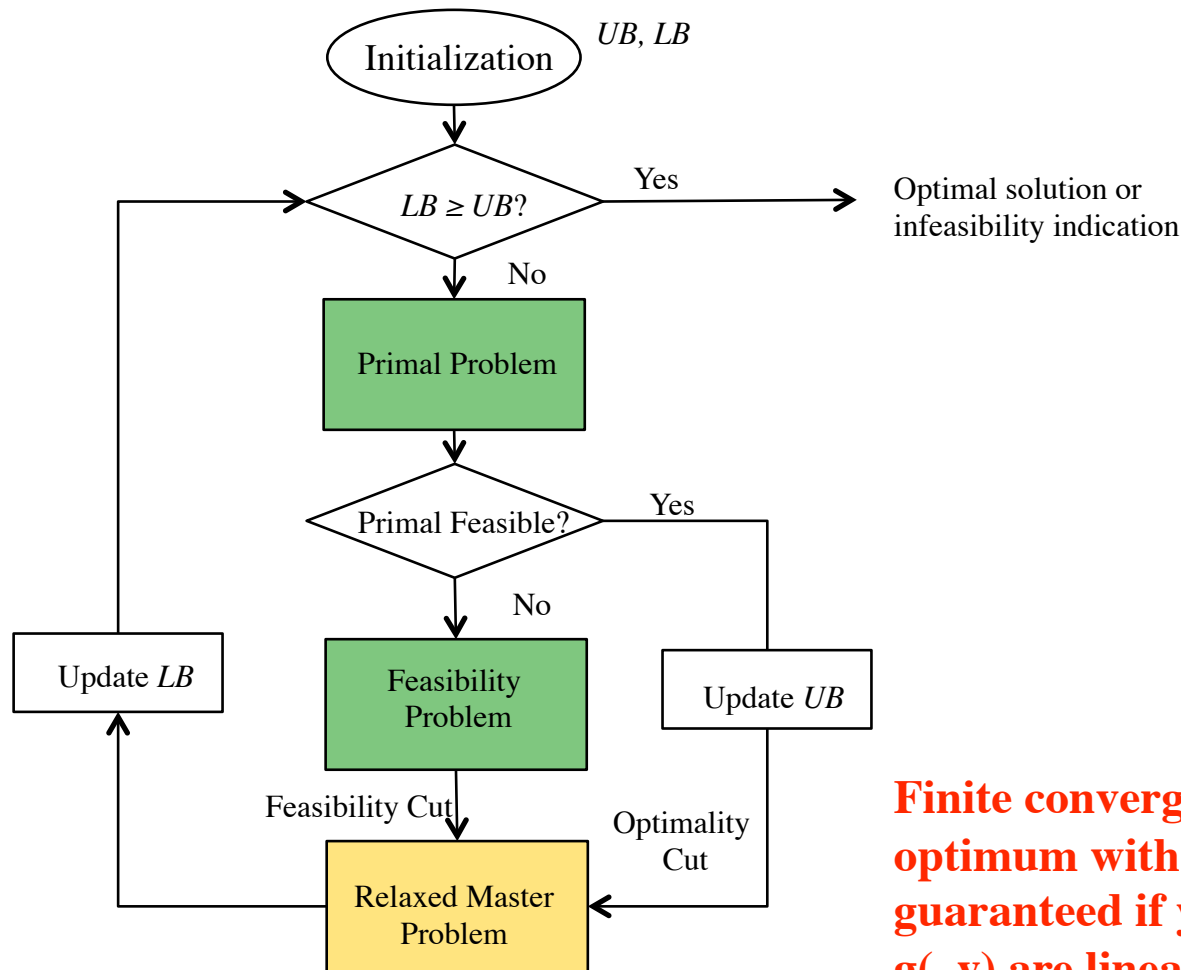
$$\begin{aligned} \min_{y, \eta} \quad & \eta \\ \text{s.t.} \quad & \eta \geq \text{obj}_{\text{Primal}}^j + f_2(y) - f_2(y^{(k)}) + (\lambda^j)^T [g_2(y) - g_2(y^{(k)})], \quad \forall \lambda^j \in T^k \\ & 0 \geq \text{obj}_{\text{Feas}}^i + (\mu^i)^T [g_2(y) - g_2(y^{(k)})], \quad \forall \mu^i \in S^k \end{aligned}$$

\* A.M. Geoffrion. Generalized Benders decomposition. Journal of Optimization Theory and Applications, 10(4):237–260, 1972.



# Generalized Benders Decomposition

## Algorithm flowchart



**Finite convergence proof see to an optimum with a given tolerance is guaranteed if  $y$  is integer or  $f(.,y)$ ,  $g(.,y)$  are linear for all  $y$  in  $Y$ .**

# Generalized Benders Decomposition

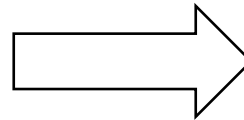
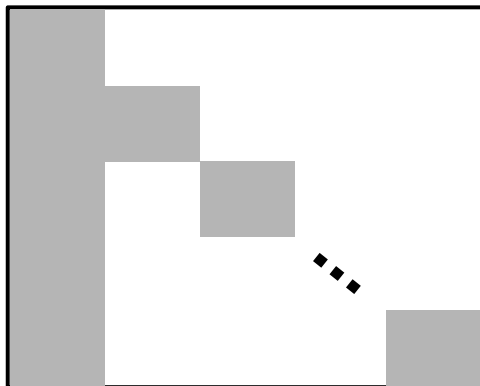
## *GBD and scenario-based stochastic programs*

$$\min_{x_1, \dots, x_s, y} \sum_{h=1}^s w_h f_h(x_h, y)$$

$$s.t. \quad g_h(x_h, y) \leq 0, \quad h = 1, \dots, s$$

$$x_h \in X_h, \quad h = 1, \dots, s$$

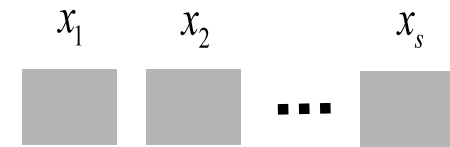
$$y \in Y$$



**Relaxed  
master  
problem**



**Primal or feasibility  
subproblems**



**But, what if convexity assumption does not hold?**

# Agenda

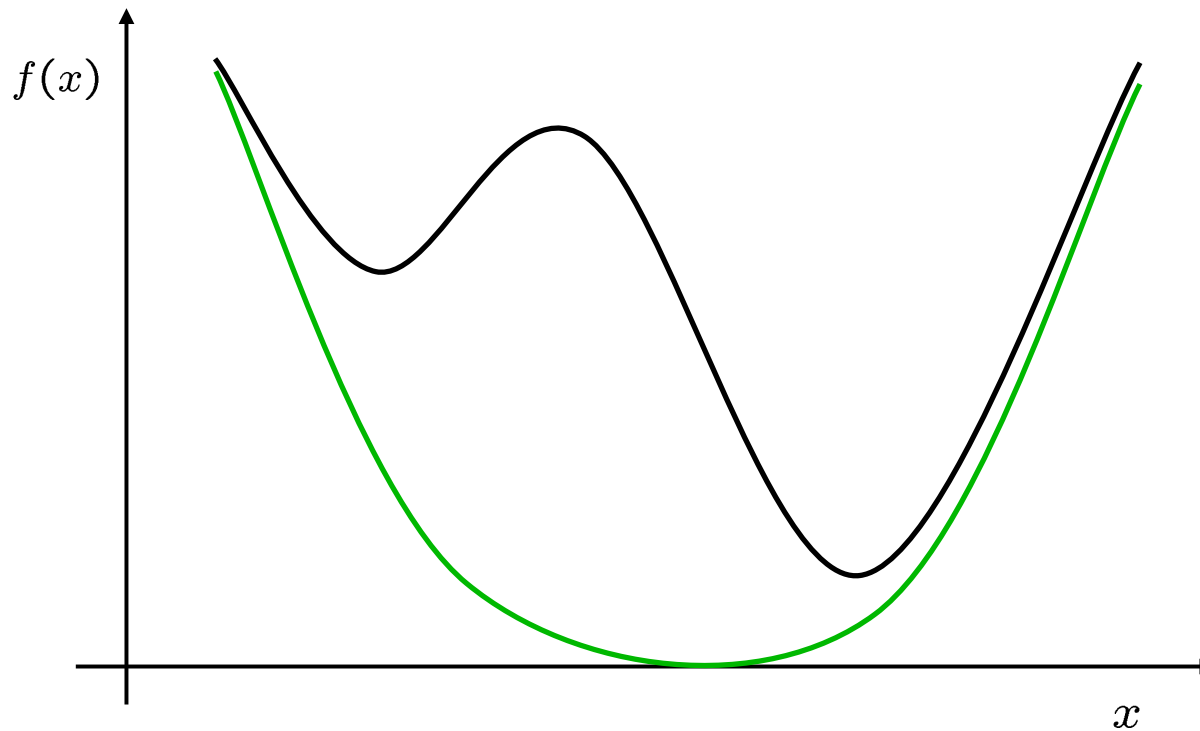
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- ◆ Motivation
- ◆ Duality Theory – A Geometric Perspective
- ◆ Generalized Benders Decomposition
- ◆ **Nonconvex Generalized Benders Decomposition**
- ◆ Computational Study
- ◆ Summary

# Nonconvex Generalized Benders Decomposition

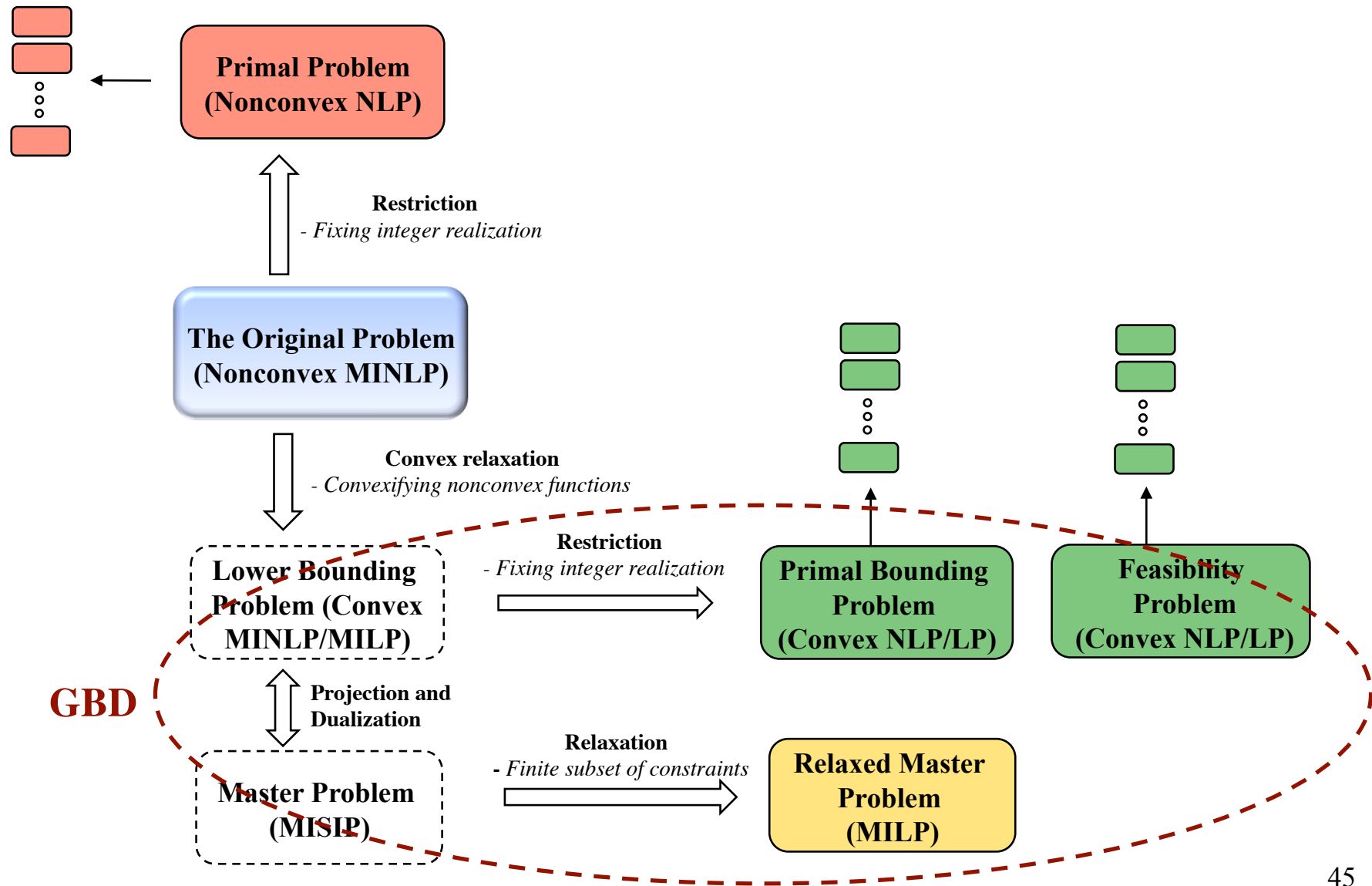
## *Convex relaxation*

---



- ◆ Convex relaxation  $\neq$  convex approximation
- ◆ Most practical problems can be relaxed via McCormick's approach
  - G. P. McCormick. Computability of global solutions to factorable nonconvex programs: Part I – Convex underestimating problems. *Mathematical Programming*, 10:147–175, 1976.
  - E. P. Gatzke, J. E. Tolsma, and P. I. Barton. Construction of convex relaxations using automated code generation technique. *Optimization and Engineering*, 3:305–326, 2002.

# Nonconvex Generalized Benders Decomposition Overview



# Nonconvex Generalized Benders Decomposition

## *Lower bounding problem via convexification*

**The Original Problem**  
(Nonconvex MINLP)

$$\begin{aligned} \min_{x_1, \dots, x_s, y} \quad & \sum_{h=1}^s w_h (f_h(x_h) + c_h^T y) \\ \text{s.t.} \quad & g_h(x_h) + B_h y \leq 0, \quad h = 1, \dots, s \\ & x_h \in X_h, \quad h = 1, \dots, s \\ & y \in Y \end{aligned}$$

**Lower Bounding Problem (Convex MINLP/MILP)**

$$\begin{aligned} \min_{x_1, \dots, x_s, y} \quad & \sum_{h=1}^s w_h (u_{f,h}(x_h, q_h) + c_h^T y) \\ \text{s.t.} \quad & u_{g,h}(x_h, q_h) + B_h y \leq 0, \quad h = 1, \dots, s \\ & x_h \in X_h, q_h \in Q_h, \quad h = 1, \dots, s \\ & y \in Y \end{aligned}$$

# Nonconvex Generalized Benders Decomposition

## *Primal bounding problem and feasibility problem*

Primal Bounding Problem (LP)



$$\begin{aligned}
 obj_{PBP_h}^{(k)} &= \min_{x_h, q_h} w_h \left( u_{f,h}(x_h, q_h) + c_h^T y^{(k)} \right) \\
 s.t. \quad & u_{g,h}(x_h, q_h) + B_h y^{(k)} \leq 0, \\
 & (x_h, q_h) \in X_h \times Q_h
 \end{aligned}$$

Feasibility Problem (LP)



$$\begin{aligned}
 obj_{FP_h}^{(k)} &= \min_{x_h, q_h, z_h} w_h \|z_h\| \\
 s.t. \quad & u_{g,h}(x_h, q_h) + B_h y^{(k)} \leq z_h, \\
 & (x_h, q_h) \in X_h \times Q_h, \\
 & z_h \in Z \subset \{z \in \mathbb{R}^m : z \geq 0\}
 \end{aligned}$$

# Nonconvex Generalized Benders Decomposition

## Master problem

Master Problem  
(MISIP)

$$\min_{\eta, y} \eta$$

$$s.t. \eta \geq \sum_{h=1}^s (w_h c_h^T + \lambda_h^T B_h) y + \sum_{h=1}^s \inf_{(x_h, q_h) \in X \times Q} [w_h u_{f,h}(x_h, q_h) + \lambda_h^T u_{g,h}(x_h, q_h)],$$

$$\forall \lambda_1, \dots, \lambda_s \geq 0,$$

$$0 \geq \sum_{h=1}^s \mu_h^T B_h y + \sum_{h=1}^s \inf_{(x_h, q_h) \in X \times Q} \mu_h^T u_{g,h}(x_h, q_h),$$

$$\forall \mu_1, \dots, \mu_s \in M,$$

$$y \in Y, \eta \in \mathbb{R}$$



# Nonconvex Generalized Benders Decomposition

## *Relaxed master problem*

---

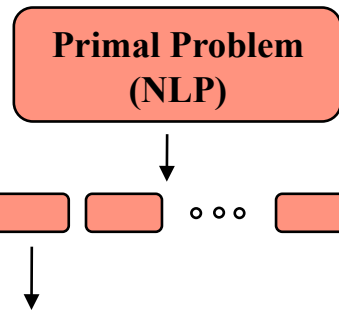
Relaxed Master  
Problem (MILP)

$$\begin{aligned}
 & \min_{\eta, y} \quad \eta \\
 & s.t. \quad \eta \geq \text{obj}_{PBP}^{(j)} + \left( \sum_{h=1}^s \left( w_h c_h^T + (\lambda_h^{(j)})^T B_h \right) \right) (y - y^{(j)}), \quad \forall j \in T^k, \\
 & \quad \quad 0 \geq \text{obj}_{FP}^{(i)} + \left( \sum_{h=1}^s (\mu_h^{(i)})^T B_h \right) (y - y^{(i)}), \quad \forall i \in S^k, \\
 & \quad \quad \sum_{r \in E^{(1)}} y_r - \sum_{r \in E^{(0)}} y_r \leq |\{r : y_r^{(t)} = 1\}| - 1, \quad \forall t \in T^k \cup S^k, \\
 & \quad \quad y \in Y, \eta \in \mathbb{R}
 \end{aligned}$$

# Nonconvex Generalized Benders Decomposition

## *Primal problem*

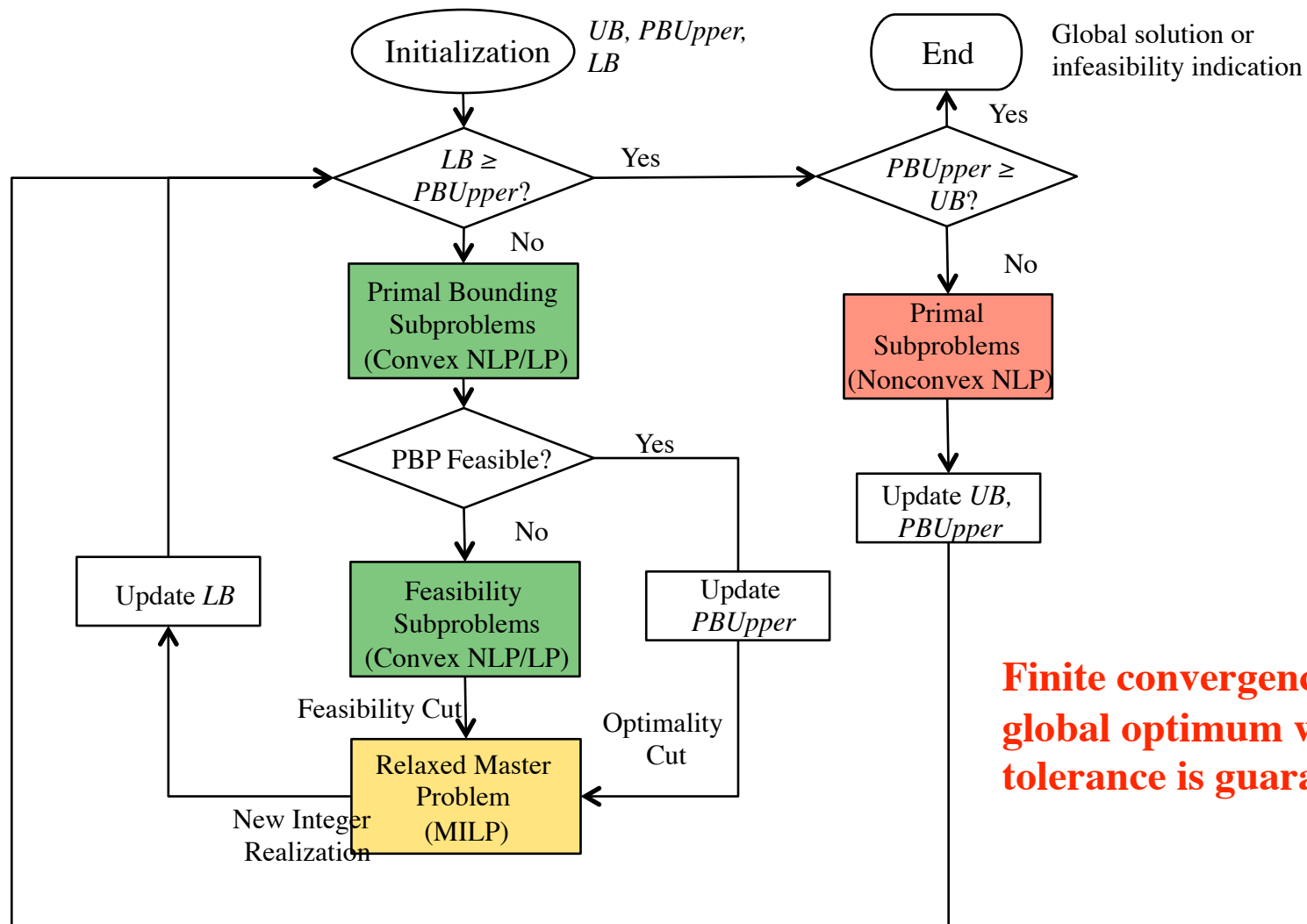
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$$\begin{aligned} \text{obj}_{PP_h}^{(l)} &= \min_{x_h} w_h \left( f_h(x_h) + c_h^T y^{(l)} \right) \\ \text{s.t. } & g_h(x_h) + B_h y^{(l)} \leq 0, \\ & f_h(x_h) + c_h^T y^{(l)} \leq UBD_h^{(l)}, \\ & x_h \in X_h \end{aligned}$$

# Nonconvex Generalized Benders Decomposition

## Algorithm flowchart



**Finite convergence to a global optimum with a given tolerance is guaranteed!**

# Agenda

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- ◆ Motivation
- ◆ Duality Theory – A Geometric Perspective
- ◆ Generalized Benders Decomposition
- ◆ Nonconvex Generalized Benders Decomposition
- ◆ **Computational Study**
- ◆ Summary

# Computational Study - Implementation Issues

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## Platform

- CPU 2.83 GHz, Memory 1 GB, Linux, GAMS 22.8.1

## Solvers

- LP and MILP solver : CPLEX
- Global NLP solver: BARON
- Local NLP solver: SNOPT

## Methods for Comparison

1. BARON – The state-of-the-art global optimization solver
2. DA – The proposed decomposition method
3. EI – Naïve integer enumeration

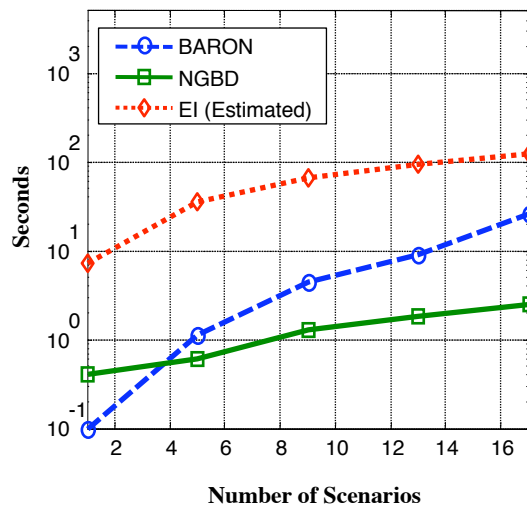
## Relative Tolerance for Global Optimization

-  $10^{-2}$

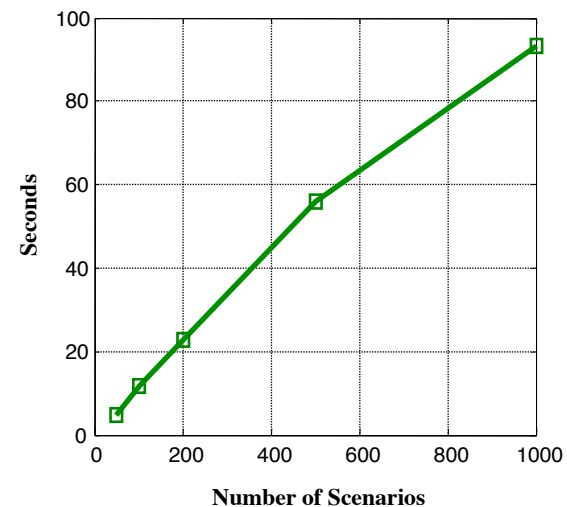
# Computational Study

## - The Stochastic Haverly Pooling Problem

The stochastic problem contains 16 binary variables and 21s continuous variables (s represents total number of scenarios).



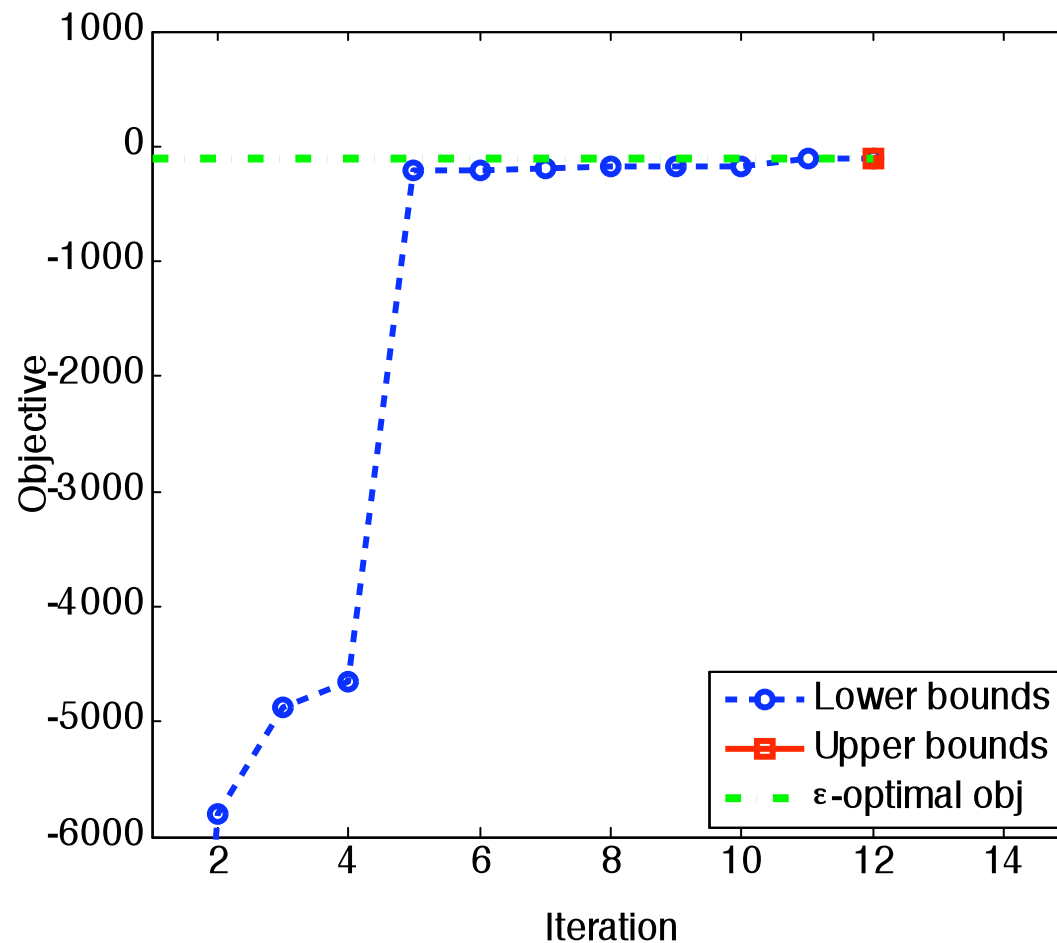
(a) Solver times with different methods



(b) Solver times with NGBD for more scenarios

# Computational Study

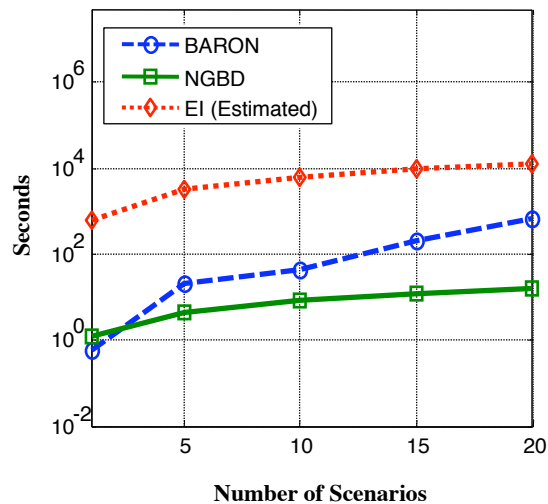
## - The Stochastic Haverly Pooling Problem



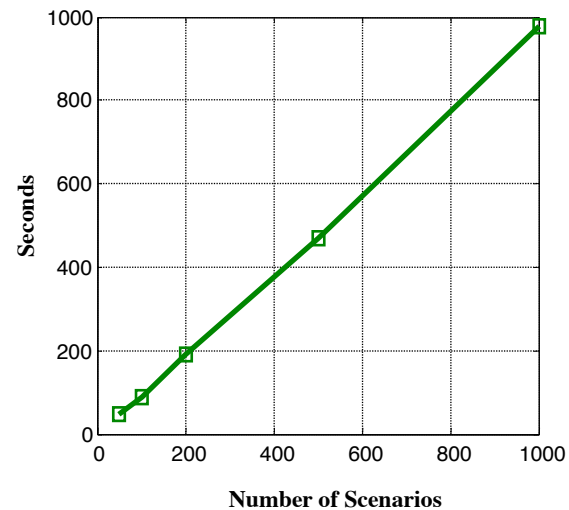
**Convergence of the upper and lower bounds over the iterations**

# Computational Study - SGPS design problem A

The stochastic problem contains 38 binary variables and 93s continuous variables (s represents total number of scenarios).



(a) Solver times with different methods

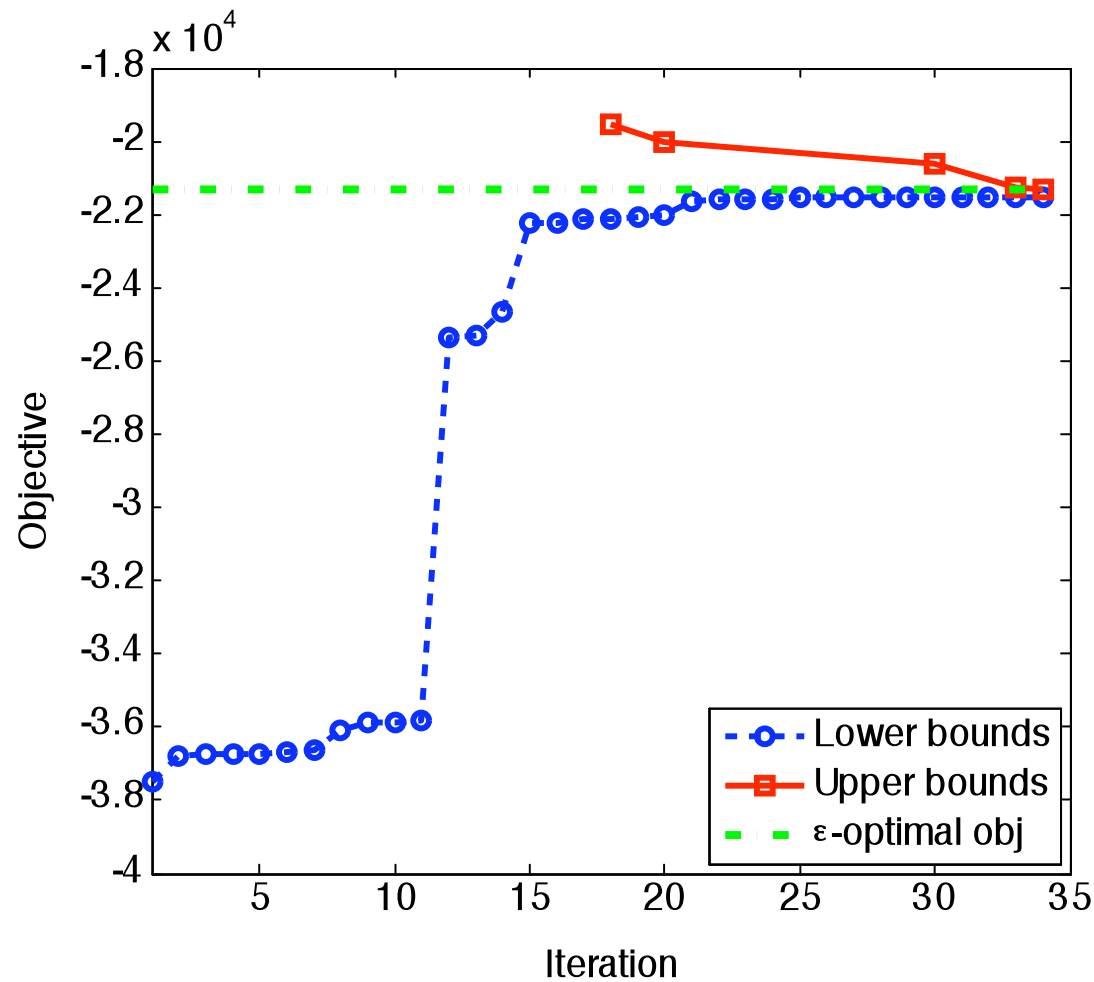


(b) Solver times with NGBD for more scenarios



# Computational Study

## - SGPS design problem A



**Convergence of the upper and lower bounds over the iterations**

# Computational Study

## - Summary of studied nonvonex problems (all solved to global optimality with given tolerances)

---

	Continuous variable /Integer variable	Time via NGBD (Second)	Nonconvexity	Implementation <sup>[1]</sup>
Haverly	21,000/16	93.4	Bilinear	(A)
Gas Network	68,000/19	15,610.7	Bilinear	(A)
SGPS A	93,000/38	978.2	Bilinear	(A)
SGPS B	93,000/38	977.1	Bilinear	(A)
SGPS C	146,410/20	4,234.8	Bilinear, quadratic, power	(B)
Software	10,648/8	260.7	Logarithmic	(B)
Pump	50,578/18	2,794.8	Bilinear, quadratic, cubic	(B)
Polygeneration	14,688/70	15,825.0 <sup>[2]</sup>	Bilinear	(C)

**Note:** [1] Problems were run with different CPUs, GAMS systems and relative termination tolerances  $\delta$ :

(A) CPU 2.83 GHz , GAMS 28.1,  $\delta=10^{-2}$  ; (B) CPU 2.83 GHz, GAMS 23.4,  $\delta=10^{-3}$ ; (C) CPU 2.66 GHz, GAMS 23.5,  $\delta=10^{-2}$  .

[2] Enhanced NGBD with tighter lower bounding problems employed.

# Agenda

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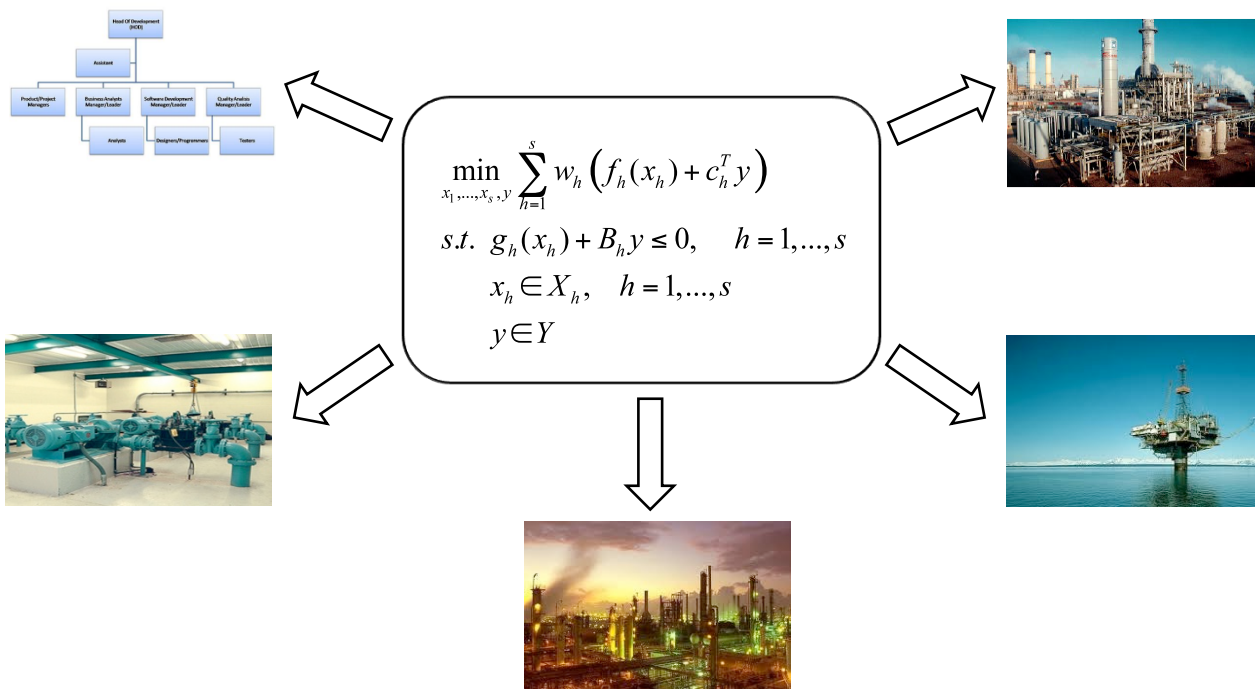
- ◆ Motivation
- ◆ Duality Theory – A Geometric Perspective
- ◆ Generalized Benders Decomposition
- ◆ Nonconvex Generalized Benders Decomposition
- ◆ Computational Study
- ◆ **Summary**

## Summary

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- ◆ Energy infrastructure design problems may be modeled as **stochastic separable mixed-integer-nonlinear programs**;
- ◆ **Generalized Benders decomposition** provides a promising nonlinear decomposition framework for scenario-based stochastic programs;
- ◆ **Nonconvex generalized Benders decomposition** tackles nonconvexity via relaxation to convex surrogate problems;
- ◆ **Nonconvex** engineering problems with practical sizes have been solved with NGBD to **global optimality** within reasonable time.

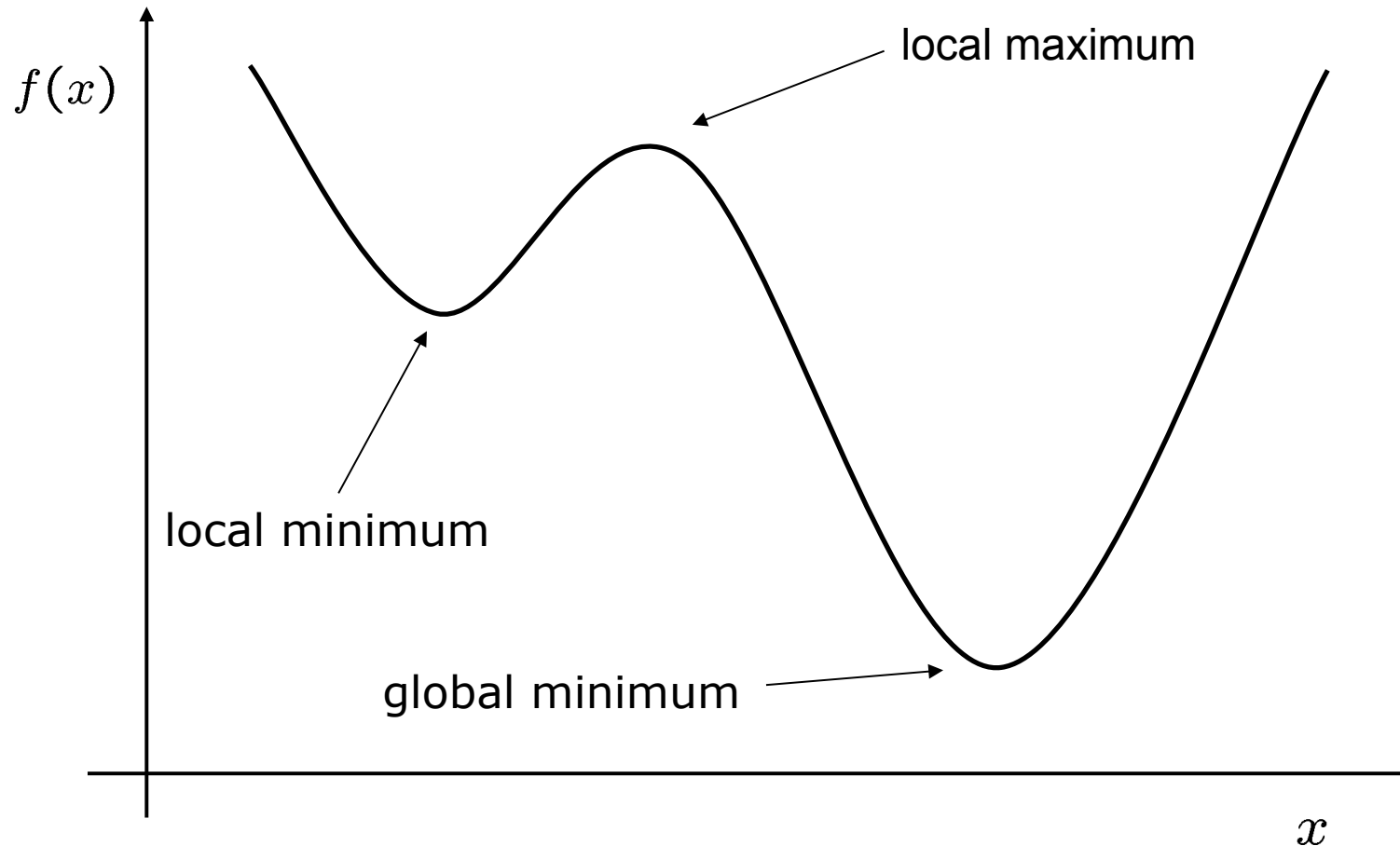
# The end, and not the end.



# Global Optimization via Branch-and-Bound

- *Nonconvex optimization*

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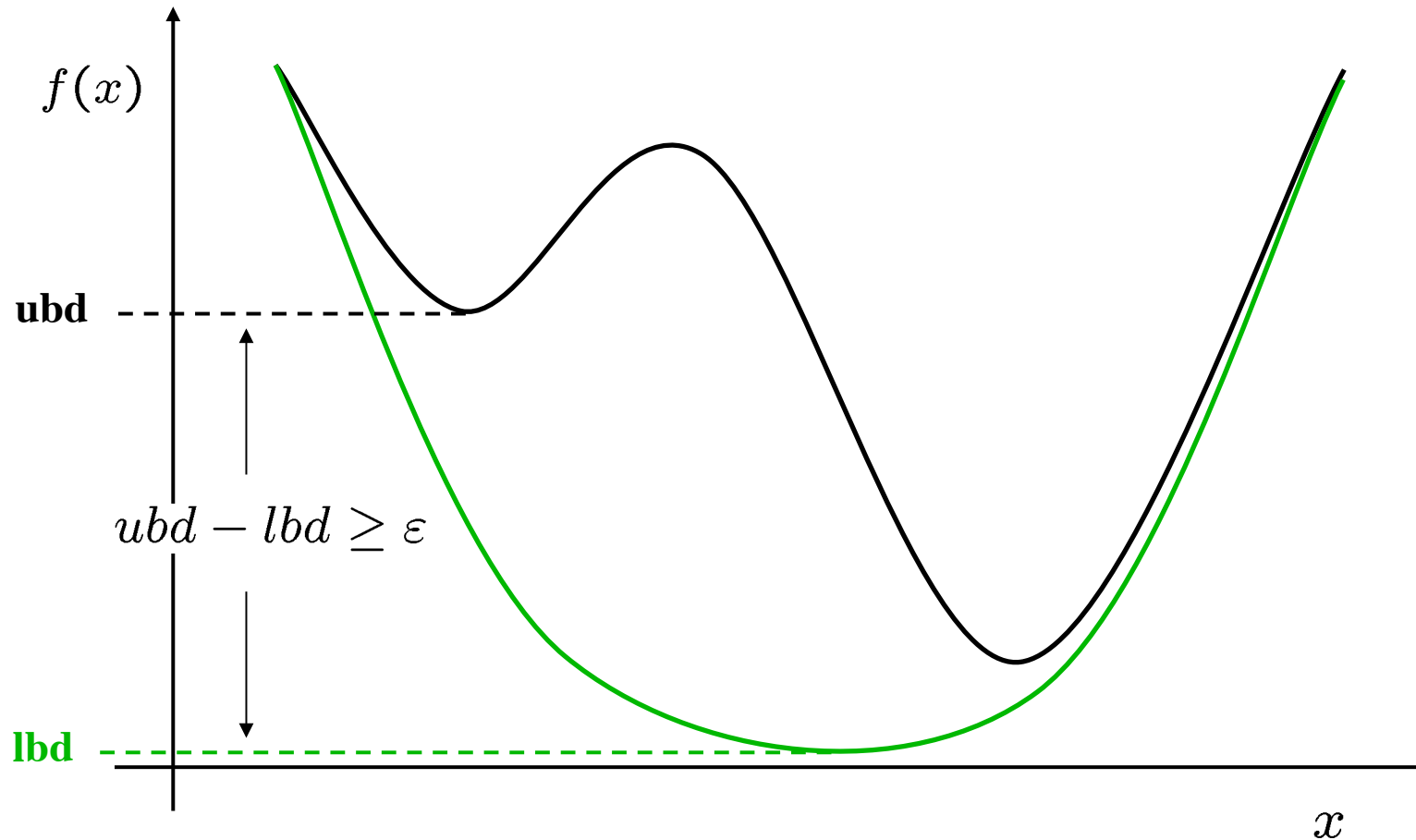


Standard optimization techniques cannot distinguish between suboptimal local minima

# Global Optimization via Branch-and-Bound

## - Convex Relaxation

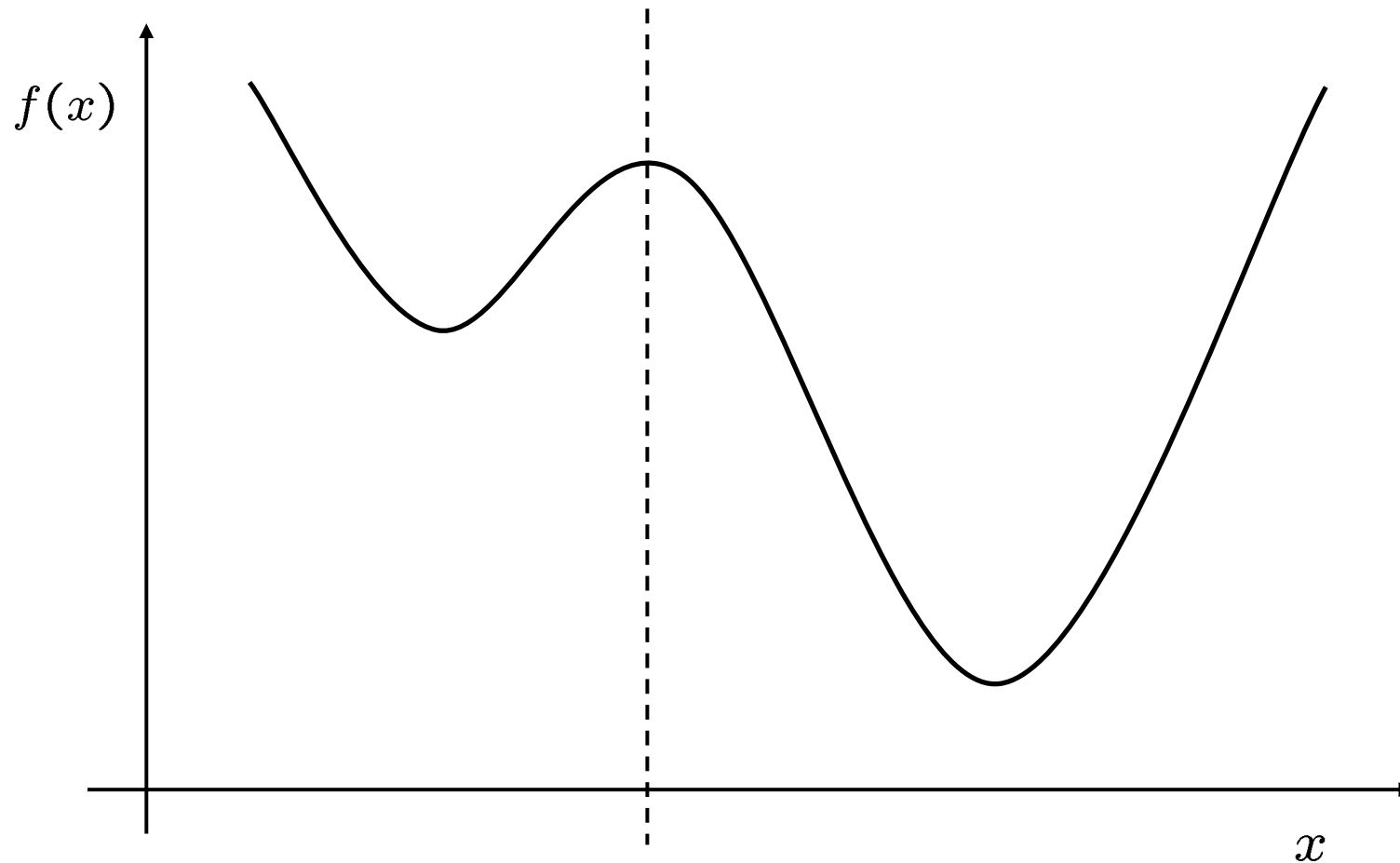
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# Global Optimization via Branch-and-Bound

## - Branch

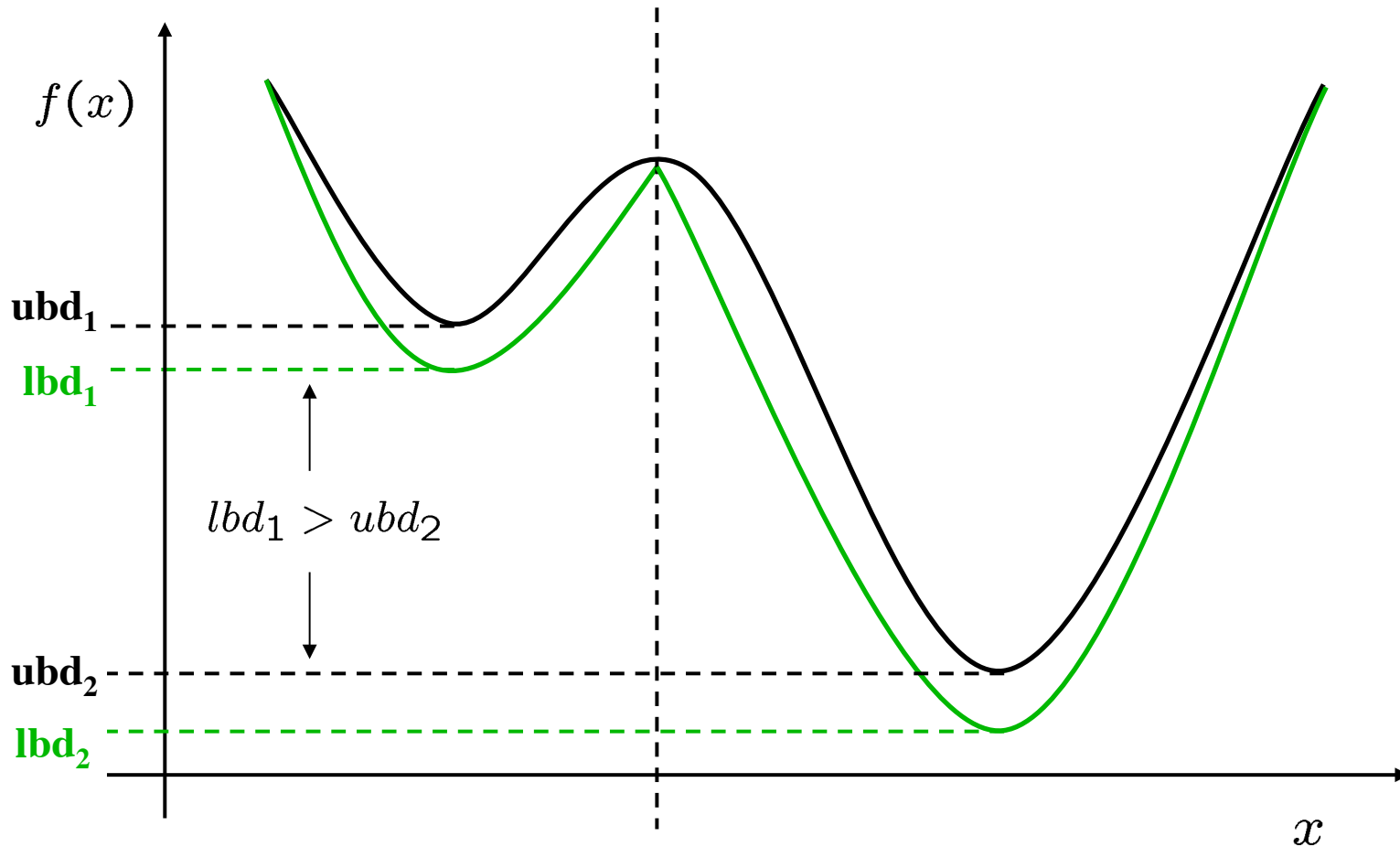
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# Global Optimization via Branch-and-Bound

- Branch, and bound



# Global Optimization via Branch-and-Bound

- Branch, bound and fathom

